

# The Jump Number of Suborders of the Power Set Order

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**Abstract.** Let  $P$  be an ordered set induced by several levels of a power set. We give a formula for the jump number of  $P$  and show that reverse lexicographic orderings of  $P$  are optimal. The proof is based on an extremal set result of Frankl and Kalai.

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## 1. Introduction

For a linear extension  $L$  of an ordered set  $P$  a pair  $(x, y)$  which is adjacent in  $L$  but incomparable in  $P$ , is called a *jump* (or *setup*). The number of jumps of  $L$  is denoted by  $s(P, L)$ , and the *jump number* of  $P$ , denoted by  $s(P)$ , is defined by  $s(P) = \min\{s(P, L) \mid L \text{ a linear extension of } P\}$ . A linear extension  $L$  of  $P$  is called *optimal* if  $s(L, P) = s(P)$ . The *jump number problem*, a special scheduling task, is to determine  $s(P)$  and to find optimal linear extensions of  $P$ .

This problem has gained a lot of attention in the last years as documented by many articles on this subject in this journal. For an introduction and references, see e.g., [1].

Let  $B_n$  denote the lattice of all subsets of an  $n$ -element set  $S$ . For a subset  $\{l_1, l_2, \dots, l_t\}$  of  $\{0, \dots, n\}$  with  $l_1 < l_2 < \dots < l_t$  we define  $B_n(l_1, \dots, l_t)$  to be the suborder of  $B_n$  which is induced by restricting  $B_n$  to the sets of cardinality  $l_1, \dots, l_t$ . We shall give a formula for the jump number of this order by proving that reverse lexicographic orderings are optimal. The proof is based on this extremal set result:

**THEOREM** (Frankl [2], Kalai [5]) *Let  $A_1, \dots, A_m$  and  $B_1, \dots, B_m$  be subsets of a set with  $|A_i| \leq a$ ,  $|B_i| \leq b$  and  $A_i \cap B_i = \emptyset$  for  $i = 1, \dots, m$ . Then  $A_i \cap B_j \neq \emptyset$  for all  $i, j \in \{1, \dots, m\}$  with  $i > j$  implies  $m \leq \binom{a+b}{a}$ .*

Given a linear extension  $L$ , we call a pair  $(x, y)$  a *bump* if  $x <_L y$  and  $x <_P y$ , i.e., if  $(x, y)$  is a covering pair in  $L$  as well as in  $P$ . The number of bumps of  $L$  is denoted by  $b(P, L)$  and  $b(P)$  is defined to be  $\max\{b(P, L) \mid L \text{ is a linear extension of } P\}$ . Obviously,  $s(P) + b(P) = |P| - 1$ .

## 2. The Result

### THEOREM.

$$s(B_n(l_1, \dots, l_t)) = -1 + \sum_{k=1}^t \binom{n}{l_k} - \sum_{k=1}^{t-1} \binom{n - l_{k+1} + l_k}{l_k}.$$

*Reverse lexicographic orderings of  $B_n(l_1, \dots, l_t)$  are optimal.*

*Proof.* We have to show that

$$b(B_n(l_1, \dots, l_t)) = \sum_{k=1}^{t-1} \binom{n - l_{k+1} + l_k}{l_k}.$$

In order to prove that the left side is less than or equal to the right side, it suffices to show that

$$b(B_n(l_k, l_{k+1})) \leq \binom{n - l_{k+1} + l_k}{l_k}.$$

Let  $L$  be a linear extension of  $B_n(l_k, l_{k+1})$  with a maximal number of bumps  $(A_1, C_1), \dots, (A_m, C_m)$ , where  $A_i, C_i \subseteq S$  and  $|A_i| = l_k$  and  $|C_i| = l_{k+1}$ . We assume also that the bumps are ordered as they occur in the linear extension  $L$ , i.e.,

$$A_1 <_L C_1 <_L A_2 <_L C_2 <_L \dots <_L A_m <_L C_m.$$

Now  $A_i \not<_L C_j$  and hence  $A_i \not\subseteq C_j$  for  $i > j$ . Setting  $B_j := S \setminus C_j$  we have  $A_i \cap B_j \neq \emptyset$  for  $i > j$  and can apply the forestanding Theorem. Thus

$$b(B_n(l_k, l_{k+1})) = m \leq \binom{n - l_{k+1} + l_k}{l_k}.$$

Now let  $S$  be ordered, say  $S = [n] := \{1, \dots, n\}$ , and let  $L$  be a reverse lexicographic ordering of  $B_n$ , i.e.,  $A <_L B$  iff  $\max((A \cup B) \setminus (A \cap B)) \in B$  for  $A, B \subseteq S$ . In order to prove that the left side of the equation above is greater or equal than the right side, it suffices to show that

$$b(B_n(l_k, l_{k+1}), L) \geq \binom{n - l_{k+1} + l_k}{l_k}.$$

We claim that all pairs  $(A, B)$  for which  $A \subseteq [n] \setminus [l_{k+1} - l_k]$  and  $B = A \cup [l_{k+1} - l_k]$ , are bumps of  $L$  in  $B_n(l_k, l_{k+1})$ . But this is clear because  $A \subseteq B$  and  $A <_L B$  and, moreover,  $A <_L B$ , since  $B_n(l_k, l_{k+1})$  is of height one. There are  $\binom{n - l_{k+1} + l_k}{l_k}$  such pairs, which finishes the proof.

COROLLARY.  $s(B_n) = 2^n - 1$ .

This can also easily be argued directly as follows. A linear extension of an ordered set  $P$  induces a chain partition  $C_1 \cup \dots \cup C_r$ . By  $l(C_i)$  we denote the length of  $C_i$ , which is the number of elements of  $C_i$  minus 1. Now  $b(P)$  equals  $\sum l(C_i)$ , where the  $C_i$  are induced by a linear extension  $L$  of  $P$ , which is chosen such that the sum is maximal. The chains have to be convex subsets of the order, which in the case of  $B_n$  implies that they are of length at most one. Now it is easy to see that  $b(B_n) = 2^n - 1$ .

In [3], Gierz and Poguntke proved that  $b(P) \leq \text{rank } M(P)$ , where  $M(P)$  denotes an incidence matrix indexed by elements of  $P$ , namely

$$(M(P))_{x,y} = \begin{cases} 1 & \text{if } x < y, \\ 0 & \text{else.} \end{cases}$$

In case of our Theorem, however, this bound does not help much, because  $M(B_n(l_k, l_{k+1}))$  has rank  $\binom{n}{l_k}$  if  $l_k + l_{k+1} \leq n$  (cf. [4]).

It should be interesting to determine the jump number of other classical ordered sets, like the partition lattice, linear lattices, and so on. Nothing seems to be known on this.

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