

# Degree-associated Reconstruction of Graphs

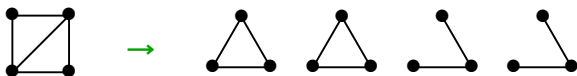
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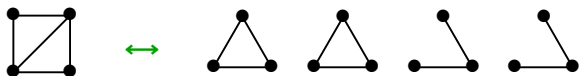
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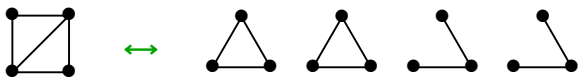


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- $rn(G) = 3$  for all trees with  $\geq 5$  vertices (Myrvold ['90]).



## Degree-associated Reconstruction

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A **dacard** ("degree-associated card") is a pair  
 $(G - v, d_G(v))$  for some  $v \in V(G)$ .

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**Ex.**  $\text{rn}(tK_p) = p + 2$  (Myrvold [1989]),  
but  $\text{drn}(tK_p) = 3$  (Ramachandran [2006]).

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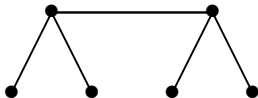
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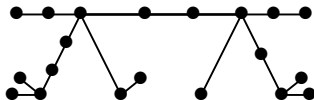
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- **Stronger:** (Spinoza [2014+]) If  $G$  arises from a caterpillar by edge-subdivisions and then growing toes (pendant edges) at leaves, then  $\text{drn}(G) \leq 2$ .

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- Among vertex-transitive graphs, always  $\text{adrn}(G) = \text{drn}(G)$ , so  $\text{adrn}(2K_{n/4, n/4}) = (n/4) + 2$ .

## Other Results on $\text{adrn}(G)$

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- Always  $\text{rn}$ ,  $\text{drn}$ , and  $\text{adrn}$  are the same for  $G$  and  $\overline{G}$ , but not so for the edge versions  $\text{ern}$ ,  $\text{dern}$ , and  $\text{adern}$ .

## Small Values

**Thm.** (Barrus–West [2010]) A single d-card  $(C, d)$  determines a graph only in the following cases.

- 1)  $d \in \{0, |V(C)|\}$ .
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**Cor.** Almost always  $rn(G) = arn(G) = 3$  and  $drn(G) = adrn(G) = 2$ .



## Regular Graphs

**Thm.** (Barrus–West [2010]) If  $G$  is  $k$ -regular with  $n$  vertices, then  $\text{drn}(G) \leq \text{adrn}(G) \leq \min\{k+2, n-k+1\}$ .

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Thus  $C = H - u$  for some  $u$ . If  $H \not\cong G$ , then  $\exists v \in N_H(u)$  such that  $d_C(v) = k$ . Thus  $d_H(v) = k+1$ .

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**Ex.** (Ramachandran [2006]). Equality holds for  $tK_{m,m}$  with  $t > 1$ ; that is,  $\text{drn}(tK_{m,m}) = m+2$ .

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**Ex.** These conditions hold for the Petersen graph, hypercubes (dimension  $\geq 3$ ), and both  $K_n \square K_2$  and  $C_n \square K_2$  (for  $n \geq 4$ ), but not for  $C_3 \square K_2$ .

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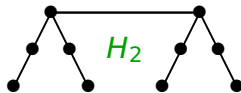
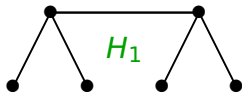
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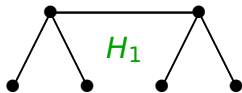


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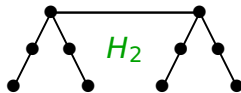
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**Thm.** (Spinoza)  $d_{rn}(T) \leq 2$  for subdivided caterpillars with toes (with few exceptions).

## Double-Brooms

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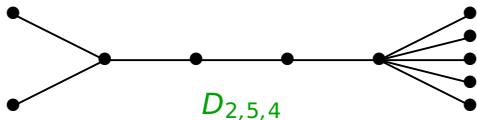
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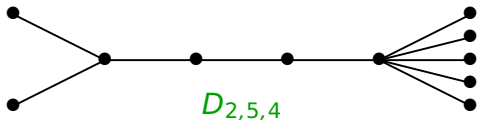
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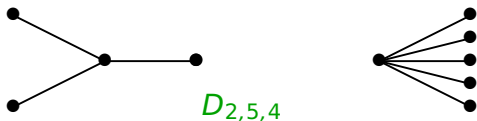
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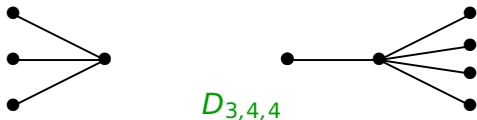
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- So far we have as large as  $\frac{1}{2}|V(T)| + 1$ .

# Open Problems

**Ques.** For vertex-transitive  $G$  with no twins, is  $\text{drn}(G)$  bounded? Always equal to 3?

**Ques.** How large can  $\text{rn}(G)$  be in terms of  $\text{drn}(G)$ ? When are they equal?

- $\text{drn}(tK_m) = 3$  (Ramachandran [2006]), but  $\text{rn}(tK_m) = m + 2$  (Myrvold [1989]).

We know no other examples with  $\text{rn}(G) > \text{drn}(G) + 1$ .

**Conj.** Only finitely many trees  $T$  satisfy  $\text{drn}(T) > 2$ .

- So far we know only  $H_1$  and  $H_2$ .

**Ques.** How large can  $\text{adrn}(T)$  be for  $n$ -vertex trees?

- So far we have as large as  $\frac{1}{2}|V(T)| + 1$ .

**Ques.** Can  $\text{rn}(G) - \text{drn}(G)$  and  $\text{adrn}(G) - \text{drn}(G)$  both be large?