

Acyclic orientations of complete bipartite graphs[☆]

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Received 2 December 1992; revised 10 August 1993

Abstract

For a complete bipartite graph, the number of dependent edges in an acyclic orientation can be any integer from $n-1$ to e , where n and e are the number of vertices and edges in the graph.

Keywords: Bipartite graph; Acyclic orientation

In combinatorics we often ask whether an integer parameter can take on all values between its extremes. In this note we consider a question of this type for acyclic orientations of a graph. An *acyclic orientation* assigns an orientation to each edge of a simple graph so that no cycle is formed.

In an acyclic orientation H of a graph G , an edge is *dependent* if reversing its orientation creates a cycle – the other edges force its orientation. This definition is due to Paul Edelman [1], who observed that the number $f(H)$ of independent edges always satisfies $n(G)-1 \leq f(H) \leq e(G)$ (where $n(G)$ and $e(G)$ denote the number of vertices and edges of G), and that these extremes are achievable when G is bipartite. Lemma 1 below includes the lower bound, and orienting all edges from one partite set to the other achieves the upper bound. Edelman asked whether G being bipartite guarantees that every number from $n(G)-1$ to $e(G)$ is achievable as $f(H)$ for some acyclic orientation H of G [3]. We call such a graph *fully orientable*. The Petersen graph, despite not being bipartite, is fully orientable, and we do not know of a triangle-free graph that is not fully orientable.

More generally, one can ask which values of $f(H)$ are achievable for an arbitrary G . It is not possible to make all three edges of a triangle independent; hence $e(G)$ may not be achievable. Indeed, the strongly connected components of an acyclic orientation of K_n must be single vertices; hence every acyclic orientation of K_n is a transitive orientation and has precisely $n-1$ independent edges. It remains open whether for

[☆] Research supported in part by NSA/MSP Grant MDA904-90-H-4011.

every graph G the achievable values of $f(H)$ form a sequence of consecutive integers beginning with $n(G) - 1$.

Independent edges are precisely those whose reversal produces another acyclic orientation. This suggests a graph $\text{AO}(G)$ on the acyclic orientations of G , in which two acyclic orientations are adjacent if one is obtained from the other by reversal of a single independent edge; the degree of an acyclic orientation in this graph is its number of independent edges. If $\text{AO}(G)$ has a Hamiltonian path, then the acyclic orientations of G can be listed in order by single-edge reversals; this condition is studied in [2].

In this note, we prove that every complete bipartite graph $K_{p,q}$ is fully orientable; Edelman [1] proved this for $p=2$. We also provide some structural lemmas about acyclic orientations.

In an arbitrary connected graph, an acyclic orientation with $f(H) = n(G) - 1$ can be constructed by using a spanning tree T found by a depth-first search from some root vertex. Orient the edges of T away from the root. Since T is a depth-first search tree, any edge not in T joins a vertex with one of its ancestors. Orient this toward the descendent; the other direction would complete a cycle. The orientation is acyclic, and the edges not in T are dependent. Lemma 1 guarantees that the edges in T are independent and that no acyclic orientation has fewer independent edges.

Lemma 1. *Every acyclic orientation of a connected simple graph G contains among its independent edges a spanning tree of G .*

Proof. Let H be an acyclic orientation of G , and let v_1, \dots, v_n be a topological ordering of the vertices of H , meaning that every edge $v_i v_j$ in H has $i < j$. Let H' be the subdigraph of H obtained by deleting all the dependent edges of H , and let G' be the underlying simple graph of H' . If G' is not connected, choose r to be the largest index such that H contains an edge $v_r v_j$ between two components of G' . Let C be the component of G' containing v_r , and choose s to be the smallest index such that $v_r v_s$ is an edge not in C . A path from v_r in H that begins along an edge in C never leaves C , by the choice of r . A path from v_r leaving C immediately by an edge other than $v_r v_s$ cannot later reach v_s , by the choice of the vertex ordering. Hence $v_r v_s$ is independent in H . The contradiction implies that G' must be connected. \square

Given an n -vertex digraph G and digraphs H_1, \dots, H_n , the *composition* $G[H_1, \dots, H_n]$ is the digraph obtained from the disjoint union $H_1 + \dots + H_n$ by adding an edge from each vertex of H_i to each vertex of H_j for each edge $v_i v_j$ in G .

Lemma 2. *If G and H_1, \dots, H_n are acyclic digraphs, then the composition $G' = G[H_1, \dots, H_n]$ is acyclic. Furthermore, if $I(G)$ denotes the set of independent edges in G and n_i, r_i, t_i , respectively, are the number of vertices, sources (indegree 0), and sinks*

(outdegree 0) in H_i , then

$$f(G') = \sum_{i=1}^n f(H_i) + \sum_{v_j v_k \in I(G)} (t_j n_k + n_j r_k - t_j r_k).$$

Proof. There is no cycle within any H_i . Since G is acyclic, no path can leave any H_i and later return to it. Hence G' is acyclic, and also the independent edges of G' within H_i are precisely the independent edges of H_i . Now consider an edge xy with $x \in H_j, y \in H_k$. If $v_j v_k$ is dependent in G , then a copy of the path making it dependent also makes xy dependent in G' . If $v_j v_k$ is independent in G , then xy is dependent if and only if x has a successor in H_j and y has a predecessor in H_k . In this case, there is an xy -path of length three making xy dependent. Otherwise, an xy -path would have to visit some H_i other than $\{H_j, H_k\}$, which would violate the independence of $v_j v_k$. Hence the number of dependent edges from H_j to H_k is $(n_j - t_j)(n_k - r_k)$, and we subtract this from $n_j n_k$ to count independent edges. \square

In an acyclic n -vertex digraph having a Hamiltonian path, the independent edges are precisely the edges of the Hamiltonian path. This enables us to construct the needed orientations for the complete bipartite graph.

Lemma 3. *Let s_1, \dots, s_m be a sequence of positive integers such that the integers with odd index sum to p and the integers with even index sum to q . Then $K_{p,q}$ has an acyclic orientation with exactly $\sum_{i=1}^{m-1} s_i s_{i+1}$ independent edges.*

Proof. Begin with a directed path with vertices v_1, \dots, v_m in order. Add an edge joining every pair of vertices having indices with opposite parity, directed toward the higher indices. This digraph G is acyclic, and by the remark its independent edges are $\{v_i v_{i+1} : 1 \leq i \leq m-1\}$. Perform the composition G' in which v_i is replaced by the digraph H_i consisting of an independent set of s_i vertices. Each vertex of H_i is both a source and a sink; $t_i = r_i = n_i = s_i$. By Lemma 2, G' is an acyclic orientation of $K_{p,q}$ with the desired number of independent edges. \square

The construction in Lemma 3 actually yields every acyclic orientation of $K_{p,q}$. To see this, let P be a maximum-length path in an acyclic orientation H of $K_{p,q}$, consisting of vertices v_1, \dots, v_m in order. The subdigraph of H induced by v_1, \dots, v_m must be the digraph G used in the construction. Let S_i be the set of vertices in H having the same predecessors and successors among v_1, \dots, v_m as v_i . The subdigraph of H induced by $\cup S_i$ is the construction G' above with $s_i = |S_i|$. It suffices to show that every vertex belongs to some S_i . An arbitrary vertex x of H is adjacent to all the vertices of even index or all the vertices of odd index in P . Also, all its predecessors in P must precede its successors, since H is acyclic. This forces it to have the same predecessors and successors in P as the unique vertex v_i between its predecessors and successors on P , and hence $x \in S_i$.

We use instances of Lemma 3 to complete the construction.

Theorem 1. For each value of k with $p+q-1 \leq k \leq pq$, the complete bipartite graph $K_{p,q}$ has an acyclic orientation with exactly k independent edges.

Proof. For $K_{1,q}$ there is nothing to prove, so we may assume $p, q \geq 2$. We need only show that all integers in the desired range are achievable in the form $f(s_1, \dots, s_m) = \sum_{i=1}^{m-1} s_i s_{i+1}$ for some positive integer sequence s_1, \dots, s_m such that the odd-indexed numbers sum to p and the even-indexed numbers sum to q .

We primarily use 6-term sequences of the form $p-1-k, 1, 1, l, k, q-l-1$, where $1 \leq k \leq p-1$ and $1 \leq l \leq q-1$. The odd terms sum to p , the even to q . All entries are non-zero, except that if $k=p-1$ or $l=q-1$, then an end-term becomes 0 and we consider instead the positive sequence with fewer terms. The value is $f = p-k+l+k(q-1)$, even when $k=p-1$ or $l=q-1$ and the sequence is shorter.

When $k=l=1$, we have $f=q+p-1$. For any fixed k , f covers a sequence of $q-1$ consecutive values as l ranges from 1 to $q-1$. The top value for k and the bottom value when k is replaced by $k+1$ are the same, since $p-k+q-1+k(q-1) = p-k-1+1+(k+1)(q-1)$. Hence there are no gaps up to the largest value achievable in this way, which occurs when $k=p-1$ and $l=q-1$ and equals $pq-p+1$.

For the remaining few values, we use sequences of the form $p-k, q-1, k, 1$, where $1 \leq k \leq p$ (again the sequence shortens when $k=p$); the value of f here is $pq-p+k$, and this completes the construction. \square

Note added in proof. In “The number of independent edges in acyclic orientations”, D.C. Fisher, K. Fraughnaugh, L. Langley, and D.B. West have proved the following results: If the chromatic number of a graph is less than its girth, then the graph is fully orientable; this includes all bipartite graphs. On the other hand, the Grötzsch graph of order 11 and chromatic number 4 has no acyclic orientation with every edge independent.

Acknowledgement

The author thanks Paul Edelman for communicating the problem, Albert Li for interesting discussions about it, and a referee for suggesting stronger versions of the Lemmas.

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