

**Midterm Math 414 (Logic), February 18, 2017**

- (1) How many words are there of length  $\leq 4$  on the alphabet  $\{a, b, c, d, e\}$ ?  
(Hint: more than 100, less than 1,000)

**Solution:** There are  $5^n$  words of length  $n$  on a five-letter alphabet, so the answer is

$$1 + 5 + 5^2 + 5^3 + 5^4 = 1 + 5 + 25 + 125 + 625 = 781.$$

- (2) Find semantically equivalent simpler propositions for:

$$p_0 \vee (p_1 \wedge p_0), \quad (p_1 \rightarrow p_0) \rightarrow p_1$$

**Solution:**  $p_0 \vee (p_1 \wedge p_0)$  is semantically equivalent to  $p_0$ , and  $(p_1 \rightarrow p_0) \rightarrow p_1$  is semantically equivalent to  $p_1$ .

- (3) Determine disjunctive normal forms for:

$$\neg(p_0 \leftrightarrow p_2), \quad (p_0 \vee p_1) \rightarrow p_2$$

**Solution:**  $(p_0 \wedge \neg p_2) \vee (\neg p_0 \wedge p_2)$  is a disjunctive normal form for  $\neg(p_0 \leftrightarrow p_2)$ , and  $(\neg p_0 \wedge \neg p_1) \vee p_2$  is a disjunctive normal form for  $(p_0 \vee p_1) \rightarrow p_2$ .

- (4) Enumerate the set of all valuations  $v$  such that for all  $n$ ,

$$v(p_n \rightarrow p_{n+1}) = 1$$

(This means in particular showing that this set is countable.)

**Solution:** One  $v$  (call it  $v_{-1}$ ) that satisfies the condition is given by  $v(p_n) = 0$  for all  $n$ . For any other  $v$  that satisfies the condition there is a least  $m$  with  $v(p_m) = 1$ , hence  $v(p_{m+1}) = 1$ , so  $v(p_{m+2}) = 1$ , and so on, that is,  $v(p_n) = 1$  for all  $n \geq m$ . Given any  $m$ , the valuation  $v_m$  defined by  $v_m(p_n) = 0$  if  $n < m$  and  $v_m(p_n) = 1$  if  $n \geq m$  is indeed a valuation satisfying the condition. So the set of valuations satisfying the condition is  $\{v_{-1}, v_0, v_1, v_2, \dots\}$ .

*For solutions of the last two problems, see back of sheet.*

- (5) Show:  $\vdash [(\phi \rightarrow \psi) \wedge (\phi \rightarrow \neg\psi)] \rightarrow \neg\phi$   
by giving a derivation.

- (6) Show:  $\vdash [(\phi \rightarrow \theta) \wedge (\psi \rightarrow \theta)] \rightarrow [(\phi \vee \psi) \rightarrow \theta]$   
by giving a derivation that uses the  $\vee$ -elimination rule at a key step, and without rewriting  $\phi \vee \psi$  as  $\neg(\neg\phi \wedge \neg\psi)$ .