

Quasirandom Hypergraphs

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Theorem

Early work by Thomason (1987) and Chung-Graham-Wilson (1989). Fix $0 < p < 1$, $\mathcal{G} = \{G_n\}_{n \rightarrow \infty}$ with $|V(G_n)| = n$ and $|E(G_n)| \geq p \binom{n}{2} + o(n^2)$. The following are equivalent.

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- **Eig:** $|\lambda_1| = pn + o(n)$ and $|\lambda_2| = o(n)$.

From k -SAT to Hypergraphs

Definition

A k -uniform hypergraph is a pair of sets $(V(H), E(H))$ where $E(H)$ is a collection of k -subsets of $V(H)$.

From a 3-SAT formula, build a 3-uniform hypergraph.

$$(x_1 \vee x_2 \vee \neg x_3) \wedge (x_1 \vee \neg x_2 \vee x_4) \wedge (\neg x_1 \vee \neg x_3 \vee x_4)$$

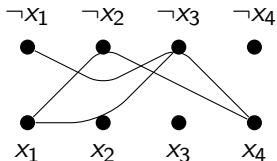
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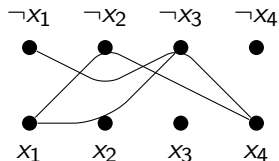
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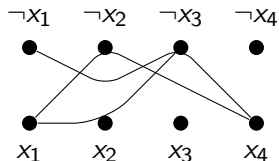


- Each truth assignment corresponds to a partition of the vertices. For example, $x_1 = T, x_2 = F, x_3 = T, x_4 = T$ corresponds to the vertex sets $T = \{x_1, \neg x_2, x_3, x_4\}$ and $F = \{\neg x_1, x_2, \neg x_3, \neg x_4\}$.

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- The assignment is a satisfying assignment if and only if F is a strong independent set.

Observation (Rödl)

For k -uniform hypergraphs with $k \geq 3$, **Disc** $\not\Rightarrow$ **Count[All]**

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Proof.

Use Erdős and Hajnal's construction: let T be a random graph tournament and form a three-uniform hypergraph by making each cyclically oriented triangle a hyperedge. There is no $K_4^{(3)}$ but **Disc** holds. □

Strong Hypergraph Quasirandomness

Theorem (Chung-Graham 1990, Kohayakawa-Rödl-Skokan 2002)

Fix $k \geq 2$, $0 < p < 1$, $\mathcal{H} = \{H_n\}_{n \rightarrow \infty}$ where H_n is k -uniform, $|V(H_n)| = n$, and $|E(H_n)| = p \binom{n}{k} + o(n^k)$. The following are equivalent.

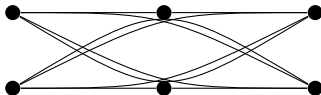
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- *CliqueDisc*[$k - 1$]
- *Count*[All]
- *Count*[Octahedron]

The octahedron is $K(\underbrace{2, \dots, 2}_k)$.



Discrepancy and Counting

Theorem (Kohayakawa-Nagle-Rödl-Schacht 2010,
Conlon-Hàn-Person-Schacht 2012)

For $k \geq 2$ and fixed $0 < p < 1$, let $\mathcal{H} = \{H_n\}_{n \rightarrow \infty}$ where H_n is k -uniform, $|V(H_n)| = n$, and $|E(H_n)| \geq p \binom{n}{k} + o(n^k)$. The following are equivalent.

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- *Disc*: $\forall U \subseteq V(H), |E(H[U])| = p \binom{|U|}{k} + o(n^k)$.
- *Expand[1 + \dots + 1]*: $\forall S_1, \dots, S_k \subseteq V(H)$,

$$e(S_1, \dots, S_k) = p \prod |S_i| + o(n^k),$$

where $e(S_1, \dots, S_k)$ is the number of tuples $(s_1, \dots, s_k) \in S_1 \times \dots \times S_k$ where $\{s_1, \dots, s_k\} \in E(H)$.

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- *Count[linear]*. A hypergraph F is linear if $\forall E_1, E_2 \in E(F)$, $|E_1 \cap E_2| \leq 1$.
- *Count[C₄]*.

Discrepancy and Counting - Definitions

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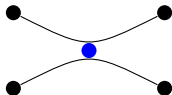
Start with a single edge.



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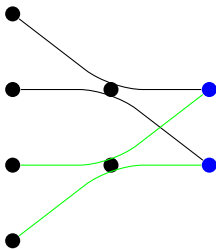
Reflect around the middle vertex.



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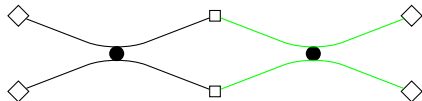
Reflect around the last two vertices.



Discrepancy and Counting - Definitions

For $k = 3$, C_4 is built as follows.

Straighten out the drawing and call it a step.

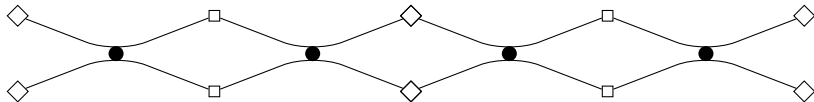


(The squares, circles, and diamonds keep track of the parts)

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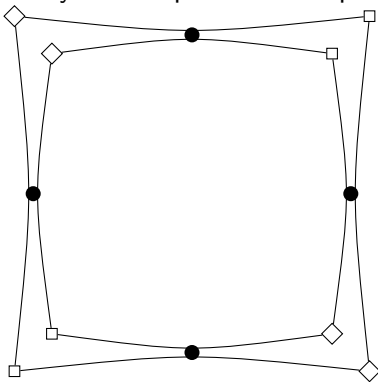
Connect steps into long paths.



Discrepancy and Counting - Definitions

For $k = 3$, C_4 is built as follows.

Identify the endpoints of the path.



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- Introduced **Eig** for hypergraphs.
- Generalized Kohayakawa-Nagle-Rödl-Schacht and Conlon-Hàn-Person-Schacht substantially.
- Determined the relationships between all studied hypergraph quasirandom properties, thereby completing a project begun by Chung in 1990.

Generalized Expansion

- For $k = 3$, **Expand[1 + 1 + 1]**: $\forall R, S, T \subseteq V(H)$,

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- **Expand[1 + 1 + 1]** \Leftrightarrow **Disc**
- **Expand[2 + 1]** \Rightarrow **Expand[1 + 1 + 1]** but the converse is false!
- More generally for $k \geq 2$, let $\pi = k_1 + \dots + k_t$ be a proper partition of k and define **Expand[π]** as follows: $\forall S_i \subseteq \binom{V(H)}{k_i}$,

$$e(S_1, \dots, S_t) = p|S_1| \cdots |S_t| + o(n^k).$$

Theorem (Lenz-Mubayi 2012+)

Fix $k \geq 2$, $\pi = k_1 + \dots + k_t$, and $0 < p < 1$. Let $\mathcal{H} = \{H_n\}_{n \rightarrow \infty}$ where H_n is k -uniform, $|V(H_n)| = n$, and $|E(H_n)| \geq p \binom{n}{k} + o(n^k)$. The following are equivalent.

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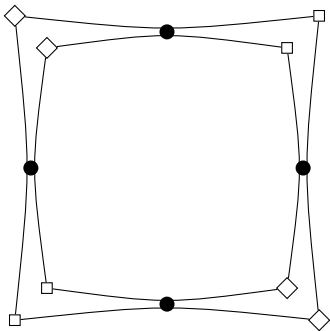
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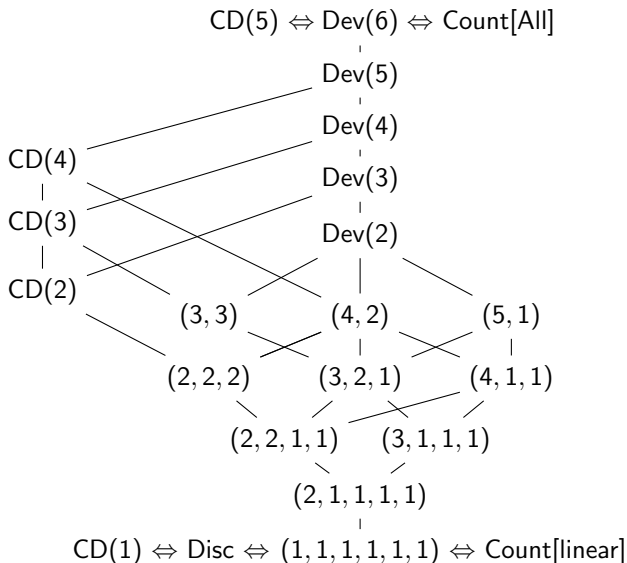
Definition

Let $k \geq 2$ and let $\pi = k_1 + \dots + k_t$ be a proper partition of k . A k -uniform hypergraph F is π -linear if there exists an ordering E_1, \dots, E_m of the edges of F such that for every i , there exists a partition of the vertices of E_i into $A_{i,1}, \dots, A_{i,t}$ such that for $1 \leq s \leq t$, $|A_{i,s}| = k_s$ and for every $j < i$, there exists an s such that $E_j \cap E_i \subseteq A_{i,s}$.

For the cycle $C_{\pi,4}$, start with C_4 defined earlier for t -uniform and expand. For $t = 3$, each diamond vertex is expanded into k_1 vertices, each square vertex into k_2 vertices, and each circle vertex into k_3 vertices.



Six Uniform Hypergraph Quasirandom Properties



Theorem (Lenz-Mubayi 2012+)

For every $k \geq 2$, $1 \leq \ell < k$, and every π , the poset of all hypergraph quasirandom properties involving *Expand* $[\pi]$, *Deviation* $[\ell]$, and *CliqueDisc* $[\ell]$ is completely determined.

Hypergraph Eigenvalues

What is $\lambda_{1,\pi}(H)$ and $\lambda_{2,\pi}(H)$? We base our definitions on definitions of Friedman and Wigderson (1995). Fix $k = 3$ and $\pi = 1 + 1 + 1$.

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- Let $W = \mathbb{R}^n$ be the vector space of dimension $n = |V(H)|$.
- The *adjacency map of H* is

$$\tau : W \times W \times W \rightarrow \mathbb{R}$$
$$\tau(e_x, e_y, e_z) = \begin{cases} 1 & \text{if } \{x, y, z\} \in E(H) \\ 0 & \text{otherwise} \end{cases}$$

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- $\lambda_{2,(1+1+1)}(H)$ is

$$\left\| \tau - \frac{k!E(H)}{n^k} J \right\|,$$

where J is the all-ones map.

Certifying Hypergraph Quasirandomness

Algorithm

On input k, ϵ, η , and H , let N be the number of labeled $C_{\pi, 4\ell}$ in H , where $\pi = \lceil k/2 \rceil + \lfloor k/2 \rfloor$ and $\ell = \lceil k/(4\epsilon) \rceil$ if k is even and $\ell = \lceil (k+1)/(4\sqrt{k}-2) \rceil$ if k is odd. Output Quasirandom if

$$N < (1 + \eta)(k!)^{4\ell} |E(H)|^{4\ell} |V(H)|^{-2\ell k}.$$

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Theorem (Lenz-Mubayi 2012+)

$\forall \epsilon \exists \eta$ so that the following holds.

- If the algorithm outputs *Quasirandom* then ***Expand*** $[\lceil k/2 \rceil + \lfloor k/2 \rfloor]$ holds.
- With high probability, the algorithm outputs *Quasirandom* on $G^{(k)}(n, p)$ if $p \geq n^{-k/2+\epsilon}$ for k even and $p \geq n^{k/2+\sqrt{k}}$ for k odd.

Refuting random k -SAT

Algorithm

On input k, ϵ, η , and a k -SAT formula ϕ , convert ϕ to a hypergraph H and run the previous algorithm counting $C_{\pi, 4\ell}$ in H . If the algorithm outputs Quasirandom, output “ ϕ is Unsatisfiable”, otherwise output Unknown.

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Theorem (Lenz-Mubayi 2012+)

$\forall \epsilon \exists \eta$ so that the following holds

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- With high probability, the algorithm outputs *Unsatisfiable* on $\text{Form}_k(n, p)$ if $p \geq n^{-k/2+\epsilon}$ for k even and $p \geq n^{k/2+\sqrt{k}}$ for k odd

Other applications and future work

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- We prove the first spectral Turán theorems for hypergraphs, proving results for cancellative hypergraphs and topologically complete hypergraphs (preprint). Currently working on F_5 -free.

Other applications and future work

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- We prove the first spectral Turán theorems for hypergraphs, proving results for cancellative hypergraphs and topologically complete hypergraphs (preprint). Currently working on F_5 -free.
- Can some graph Ramsey results that rely on graph eigenvalues be translated to hypergraphs now that we know $\text{Eig}[\pi]$?