

A combinatorial approach to the study of divisors on $\overline{M}_{0,n}$

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Goal

Goal: Illustrate how a problem from algebraic geometry can be approached using combinatorics

- 1 The combinatorial Problem
 - The space
 - The players
 - The game
- 2 The cones \mathfrak{U} and \mathfrak{L} for the space of phylogenetic trees
 - The cone \mathfrak{U}
 - The cone \mathfrak{L}
- 3 The algebraic geometry story
 - Moduli spaces
 - Divisors
 - Useful tool

Outline

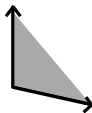
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Cones

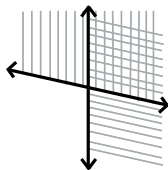
Definition

A **cone** is the positive span of a finite number of vectors, i.e., a set of the form

$$\text{pos}(v_1, \dots, v_k) := \{\lambda_1 v_1 + \dots + \lambda_k v_k : \lambda_i \geq 0\}$$



Cones can also be expressed as a finite intersection of halfspaces.

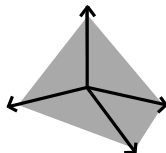


Fans

Definition

A **fan** is a family of nonempty cones such that

- 1 Every nonempty face of a cone in the fan is also a cone of the fan,
- 2 the intersection of any two cones is a face of both.



Important example: Space of Phylogenetic trees

Definition

- A **rooted tree** is a graph that has no cycles and which has a vertex of degree at least 2 labelled as the root of the tree.
- The leaves of the tree are all the vertices of degree 1; we label them from 1 to n .

Each vertex of the tree corresponds to a subset of $\{1, \dots, n\}$

Example

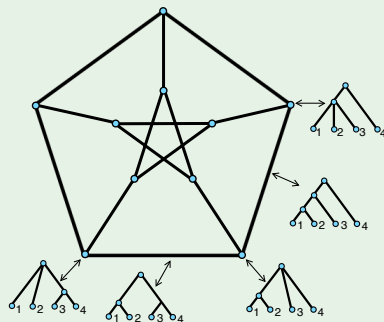


Coarse subdivision on $\mathcal{T}(K_n)$

- There is a fan whose cones are in 1-1 correspondence with rooted trees with n labelled leaves.
- Maximal cones correspond to binary trees.
- Rays correspond to subsets of $\{1, \dots, n\}$ of size ≥ 2 , so a cone corresponding to the tree T is generated by the rays corresponding to the vertices of T .
- The union of the cones of this fan is the space of phylogenetic trees

Example

$n=4$



Star¹-convex functions

Definition

Given a fan Δ , $N(\Delta)$ is the set of piecewise linear functions $\varphi : \bigcup_{\sigma \in \Delta} \sigma \rightarrow \mathbb{R}$ that are linear on each cone of Δ .

$N(\Delta)$ is isomorphic to $\mathbb{R}^{\# \text{ of rays}}$, i.e., a function φ is determined by its values on the rays.

Phylogenetic case

A function $\varphi \in N(\Delta)$ is determined by the values on the rays v_l where l is a subset of $\{1, \dots, n\}$ of size ≥ 2 .

Definition

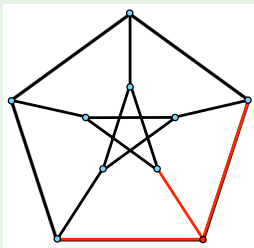
Let $\sigma \in \Delta$, we say that $\varphi \in N(\Delta)$ is **star¹-convex** on σ if it satisfies that

$$\varphi(u_1 + \cdots + u_k) \leq \varphi(u_1) + \cdots + \varphi(u_k)$$

for each u_1, \dots, u_k such that

- 1 $u_1 + \cdots + u_k \in \sigma$, and
- 2 each $u_i \in \tau_i$ where $\tau_i \supset \sigma$ and $\dim(\tau_i) = \dim(\sigma) + 1$, i.e., $\tau_i \in \text{star}^1(\sigma)$.

Example



$\text{star}^1(\sigma)$ of the cone σ corresponding to the red vertex.

Cones on $N(\Delta)$

Definition

Let σ be a cone of a fan Δ , define

- $\mathfrak{C}(\sigma)$, the set of functions $\varphi \in N(\Delta)$ that are star^1 -convex on σ ,
- the set of functions in $N(\Delta)$ star^1 convex on all cones $\sigma \in \Delta$:

$$\mathfrak{L}(\Delta) := \bigcap_{\sigma \in \Delta} \mathfrak{C}(\sigma), \text{ and}$$

- the set of functions in $N(\Delta)$ star^1 convex on all cones $\sigma \in \Delta$ of codimension 1:

$$\mathfrak{L}(\Delta) := \bigcap_{\sigma \in \Delta, \text{codim}(\sigma)=1} \mathfrak{C}(\sigma).$$

Question

Clearly $\mathfrak{L}(\Delta) \subseteq \mathfrak{L}(\Delta)$, but are the two cones equal for certain fans?

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The cone \mathfrak{U}

Theorem

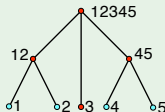
Trees corresponding to cones of codimension 1 have only one vertex with exactly 3 children. Each of these trees gives a halfspace for \mathfrak{U} which depends only on this vertex and its children.

Example

For K_5 , \mathfrak{U} is the intersection of 65 halfspaces in \mathbb{R}^{26} . Some of the halfspaces:



$$\varphi(123) + \varphi(4) + \varphi(5) + \varphi(12345) \leq \varphi(1234) + \varphi(1235) + \varphi(45)$$



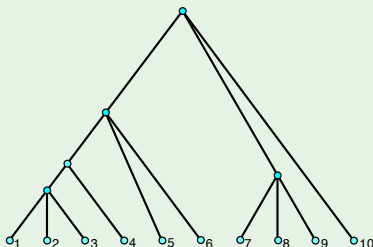
$$\varphi(12) + \varphi(3) + \varphi(45) + \varphi(12345) \leq \varphi(123) + \varphi(345) + \varphi(1245)$$

The cone \mathfrak{L}

Theorem

Let T be the tree corresponding to the cone σ and r_1, \dots, r_k be all the vertices of T having ≥ 3 children. Then computing cone $\mathfrak{C}(\sigma)$ can be reduced to computing smaller cones $\mathfrak{C}(r_i)$ where each such cone depends only on vertex r_i and its children.

Example

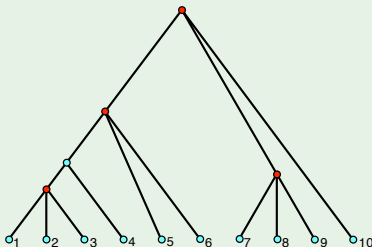


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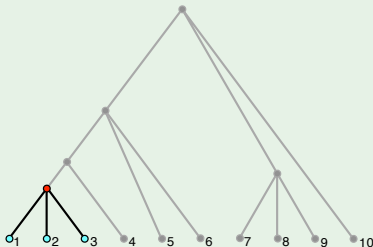


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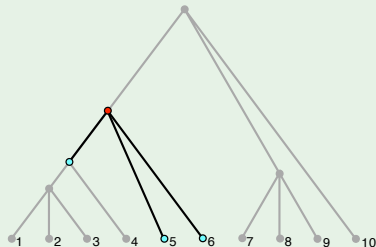
$$\varphi(1) + \varphi(2) + \varphi(3) + \varphi(123) \leq \varphi(12) + \varphi(13) + \varphi(23)$$

The cone \mathfrak{L}

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Example



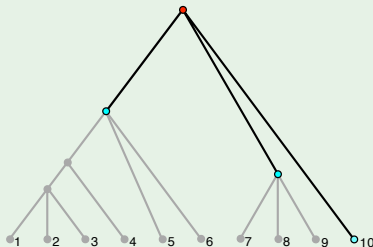
$$\varphi(1234) + \varphi(5) + \varphi(6) + \varphi(123456) \leq \varphi(12345) + \varphi(12346) + \varphi(56)$$

The cone \mathfrak{L}

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Example



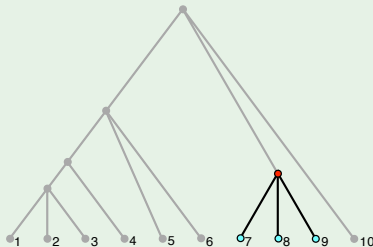
$$\varphi(123456) + \varphi(789) + \varphi(10) + \varphi(12345678910) \leq \\ \varphi(123456789) + \varphi(12345610) + \varphi(78910)$$

The cone \mathfrak{L}

Theorem

Let T be the tree corresponding to the cone σ and r_1, \dots, r_k be all the vertices of T having ≥ 3 children. Then computing cone $\mathfrak{C}(\sigma)$ can be reduced to computing smaller cones $\mathfrak{C}(r_i)$ where each such cone depends only on vertex r_i and its children.

Example



$$\varphi(7) + \varphi(8) + \varphi(9) + \varphi(789) \leq \varphi(78) + \varphi(79) + \varphi(89)$$

Inductive approach

Theorem

If $\mathfrak{L} = \mathfrak{L}$ for the space of phylogenetic trees with $n - 1$ leaves and \mathfrak{L} is contained in the intersection of the cones given by the trees with only one internal vertex and at most n leaves, then $\mathfrak{L} = \mathfrak{L}$ for the space of phylogenetic trees with n leaves.

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General idea

Theorem (Gibney and Maclagan)

The cones \mathfrak{L} and \mathfrak{U} give us a tool to compute an important cone which arises in algebraic geometry.

- A central goal in algebraic geometry is to understand maps $X \rightarrow \mathbb{P}^k$, for a projective variety X .
- A main tool in studying these maps is the nef cone of X .
- Interesting unknown case when $X = \overline{M}_{0,n}$.

Moduli spaces

Definition

The moduli space $M_{0,n}$ is a geometric space whose points correspond to isomorphism classes of smooth curves of genus 0 with n distinct marked points.

- $M_{0,n} = \{\mathbb{P}^1 \text{ with } n \text{ distinct marked points}\} / \text{automorphisms.}$
- Smooth space of dimension $n - 3$.
- Understanding of this space tells us a lot about curves.

Deligne-Mumford compactification $\overline{M}_{0,n}$

Add every curve with n marked points whose group of automorphisms fixing those points is finite.

Divisors (simplified)

Definition

A **Divisor** of a variety X is a finite sum of the form $\sum_i a_i D_i$ where each $a_i \in \mathbb{R}$ and each D_i is a codimension 1 subvariety of X .

Definition

- Let D be a divisor and C a curve in X , then $D \cdot C := \sum_i a_i |D_i \cap C|$.
- The **nef cone** of X is the cone generated by divisors such that $D \cdot C \geq 0$ for all curves C .

Example

$D = \sum_i a_i D_i$ with $a_i \geq 0$ is in the nef cone.

Useful tool

There is a natural embedding of $\overline{M}_{0,n} \hookrightarrow X_{\Delta}$, where X_{Δ} is the toric variety of the fan of the space of phylogenetic trees.

The cones $\mathfrak{L}(\Delta)$ and $\mathfrak{U}(\Delta)$ are cones of divisors on X_{Δ} .

Gibney and Maclagan use this embedding to pull back these cones to cones of divisors on $\overline{M}_{0,n}$ which give upper and lower bounds for the nef cone of $\overline{M}_{0,n}$.

If we can prove $\mathfrak{L}(\Delta_n) = \mathfrak{U}(\Delta_n)$, where Δ_n is the space of phylogenetic trees with n leaves, then we would have a nice description of the nef cone of $\overline{M}_{0,n}$, which is in general hard to compute.

This technique can also be applied to other projective varieties X for which there is a nice embedding to a toric variety.

Thank you!