

## 1 Lecture 5: The number of Sidon sets (1/25/2019)

A set  $A \subseteq [n]$  is a *Sidon set* if there is no *Sidon 4-tuple*, i.e.,  $(a, b, c, d)$  in  $A$  with  $a + b = c + d$  and  $\{a, b\} \cap \{c, d\} = \emptyset$ . A major problem concerning Sidon sets is to find the largest cardinality of a Sidon set contained in  $[n]$ . A trivial upper bound,  $|A| \leq 2\sqrt{n}$ , can be given by observing that in a Sidon set  $A$ , the sums of two different numbers in  $A$  are all distinct, i.e.  $\binom{|A|}{2} \leq 2n$ . By the works of Erdős and Turán [6], Singer [10], Erdős [4], and Chowla [3], it is known that the maximum Sidon set is of size  $(1 + o(1))\sqrt{n}$ . This is also true for the Sidon sets in  $\mathbb{Z}_n$ .

### 1.1 A connection with $C_4$ -free graphs

Due to Brown [1] and Erdős, Rényi, and Sós [5], it is well known that the maximum number of edges in a  $C_4$ -free graph is  $(\frac{1}{2} + o(1))n^{3/2}$ . The upper bound is relatively easy to obtain. However, it is difficult to construct  $n$ -vertex graphs without 4-cycles that have  $(\frac{1}{2} + o(1))n^{3/2}$  edges. Next, we show a construction for the lower bound.

**Construction 1.1** (The incidence graph of a finite projective plane, Reiman [8]). Let  $q$  be a prime power. It is known that there exists a finite projective plane with a point set  $\mathcal{P}$  and a line set  $\mathcal{L}$  such that the following hold.

- $|\mathcal{P}| = |\mathcal{L}| = q^2 + q + 1$ ;
- Each line contains exactly  $q + 1$  points;
- Each point is contained in exactly  $q + 1$  lines.

Let  $G$  be the graph with vertex set  $\mathcal{P} \cup \mathcal{L}$  and edge set

$$E = \{(p, \ell) : p \in \mathcal{P}, \ell \in \mathcal{L}, p \text{ is contained in the line } \ell\}.$$

The graph  $G$  is a bipartite graph and also an optimal  $C_4$ -free graph. Why is  $G$   $C_4$ -free? Suppose that  $G$  contains a 4-cycle  $p_1\ell_1p_2\ell_2$ . Then we would have two different points  $p_1, p_2$  and two different lines  $\ell_1, \ell_2$  such that  $p_1, p_2 \in \ell_1, \ell_2$ , which is impossible.

To construct such a projective plane, one can consider equivalence classes of the finite field  $\mathbb{F}_q^3$ ; two elements  $(a, b, c)$  and  $(a', b', c')$  in  $\mathbb{F}_q^3 \setminus \{(0, 0, 0)\}$  are equivalent if  $(a, b, c) = \lambda(a', b', c')$  for some  $\lambda \neq 0$ . Let the point set  $\mathcal{P}$  be the set of equivalence classes. There are  $\frac{q^3-1}{q-1} = q^2 + q + 1$  such points. For every  $(a, b, c) \in \mathcal{P}$ , a straight line is defined as the set of all points  $(x, y, z)$  which satisfy the equation  $ax + by + cz = 0$ . We let  $\mathcal{L}$  be the set of all such lines.

To obtain a non-bipartite construction, one can consider the projective norm graph, which was used in [5]. Let  $G$  be a graph with vertex set  $\mathcal{P}$ , where  $\mathcal{P}$  is the set of equivalence classes of  $\mathbb{F}_q^3$  defined as above and  $|\mathcal{P}| = q^2 + q + 1$ . The equivalence class of  $(a, b, c)$  is connected by an edge to the class of  $(x, y, z)$  if  $(a, b, c)$  and  $(x, y, z)$  are in two different classes, and  $ax + by + cz = 0$ . One can show that  $G$  is  $C_4$ -free with  $e(G) = \frac{1}{2}q(q+1)^2$ , and we omit the details of the proof here.

Now let  $A \subset [n]$  be a Sidon set. We define a bipartite graph  $G$  on the vertex set  $[n] \cup [2n]$  such that for two vertices  $i \in [n]$  and  $j \in [2n]$ ,  $i$  is adjacent to  $j$  if and only if there exists a number  $a \in A$  such that  $i + a = j$ . If  $A$  is of size  $C\sqrt{n}$ , then we have  $e(G) = n \cdot |A| = Cn^{3/2}$ . We claim that  $G$  is a  $C_4$ -free graph. Suppose that there exists a 4-cycle  $i, i+a, j, j+d$ , where  $a, d \in A$ . Since  $(j, i+a)$  and  $(i, j+d)$  are edges, we can find two numbers  $b, c \in A$  such that  $j + c = i + a$  and  $i + b = j + d$ . However, this also gives  $a + d = b + c$ , where  $\{a, d\} \neq \{b, c\}$  because of  $i \neq j$ . Therefore, we find a Sidon 4-tuple  $(a, d, b, c)$  in  $A$ , which is a contradiction.

## 1.2 The number of Sidon sets

In 1990, Cameron and Erdős [2] proposed the problem of determining the number of Sidon sets in  $[n]$ . By the largest cardinality of a Sidon set, it is trivial to obtain the following bound:

$$2^{(1+o(1))\sqrt{n}} \leq \text{the number of Sidon sets} \leq \binom{n}{(1+o(1))\sqrt{n}} \leq 2^{\sqrt{n} \log n}.$$

It seems natural to think that the trivial lower bound is the best possible. However, surprisingly, Saxton and Thomason [9] showed that the lower bound is not tight.

**Theorem 1.2** (Saxton-Thomason [9]). *The number of Sidon sets in  $[n]$  is at least  $2^{1.16\sqrt{n}}$ .*

*Proof.* Assume that  $n$  is a multiple of 4. Let  $A \subseteq [n/4]$  be the largest Sidon set in  $\mathbb{Z}_{n/4}$ , and  $|A| = (\frac{1}{2} + o(1))\sqrt{n}$ . For every  $a \in A$ , we can choose  $u_a$  to be any one of  $\{a, \frac{n}{4} + a, \frac{n}{2} + a, \frac{3}{4}n + a\}$ , or nothing. Let  $U = \{u_a : a \in A\}$ . Clearly, the number of choices for  $U$  is  $5^{\frac{1}{2}\sqrt{n}} \approx 2^{1.16\sqrt{n}}$ . We claim that  $U$  is a Sidon set in  $[n]$ . Suppose that there is a Sidon 4-tuple

$(a', b', c', d')$  in  $U$  such that  $a' + b' = c' + d'$ . Then we can find numbers  $a, b, c, d \in A$  such that  $a + b \equiv c + d \pmod{n/4}$ , which contradicts the fact that  $A$  is a Sidon set in  $\mathbb{Z}_{n/4}$ .  $\square$

The best result on the upper bound is given by Kohayakawa, Lee, Rödl and Samotij [7], using the graph container method.

**Theorem 1.3** (Kohayakawa-Lee-Rödl-Samotij [7]). *The number of Sidon sets in  $[n]$  is at most  $2^{c\sqrt{n}}$ , where  $c$  is a constant arbitrarily close to  $\log_2(32e) \approx 6.442$  for sufficiently large enough  $n$ .*

We remark that Saxton and Thomason [9] obtained the upper bound  $2^{(55+o(1))\sqrt{n}}$  by applying the hypergraph container method.

### 1.3 Supersaturation

For two sets  $A, U \subseteq [n]$ , define a multigraph  $H^U(A)$  on vertex set  $A$  such that for every  $a_1, a_2 \in A$  with  $a_1 < a_2$ , the multiplicity of the edge  $a_1 a_2$  in  $H^U(A)$  is the number of ordered pairs  $(u_1, u_2)$  in  $U$  such that  $(a_1, u_1, u_2, a_2)$  is a Sidon 4-tuple.

**Lemma 1.4.** *Let  $A, U \subseteq [n]$ . If  $|A| \cdot |U| \geq 6n$ , then  $e(H^U(A)) > \frac{|A|^2|U|^2}{12n}$ .*

*Proof.* Let  $F$  be a simple bipartite graph defined on the set  $A \cup [2n]$  satisfying that for every  $a \in A$  and  $m \in [2n]$ ,  $a$  is adjacent to  $m$  if and only if there is an element  $u \in U$  such that  $a + u = m$ . Clearly, for every vertex  $a \in A$ , we have  $d_F(a) = |U|$ .

Let  $\mathcal{P}$  be the set of paths of length 2 (or 3-paths) in  $F$  with endpoints in  $A$ . Then we have

$$|\mathcal{P}| = \sum_{m \in [2n]} \binom{d_F(m)}{2} \geq 2n \binom{\frac{\sum_{m \in [2n]} d_F(m)}{2n}}{2} = 2n \binom{\frac{|A| \cdot |U|}{2n}}{2} > \frac{|A|^2|U|^2}{6n}.$$

A path  $P = \{xyz\} \in \mathcal{P}$  is called *trivial* if  $x + z = y$ ; otherwise,  $P$  is *nontrivial*. Note that  $P$  is trivial if and only if both  $x$  and  $z$  belong to  $A \cap U$ . Thus, the number of trivial paths in  $\mathcal{P}$  is exactly  $\binom{|A \cap U|}{2}$ . Let  $\mathcal{P}'$  be the set of nontrivial paths in  $\mathcal{P}$ . Every 3-path in  $\mathcal{P}'$  corresponds to an edge in  $H^U(A)$  and vice versa. Therefore, we obtain

$$e(H^U(A)) = |\mathcal{P}'| = |\mathcal{P}| - \binom{|A \cap U|}{2} > \frac{|A|^2|U|^2}{6n} - \frac{|A| \cdot |U|}{2} \geq \frac{|A|^2|U|^2}{12n},$$

where the first inequality is given by  $|A \cap U| \leq \min\{|A|, |U|\} \leq \sqrt{|A| \cdot |U|}$  and the second inequality follows from the assumption  $|A| \cdot |U| \geq 6n$ .  $\square$

## 1.4 Proof of Theorem 1.3

**Proposition 1.5.** *Let  $U \subseteq [n]$  be a Sidon set, and  $A \subseteq [n] - U$ . Then  $H^U(A)$  is a simple graph.*

*Proof.* Suppose that  $a_1 a_2$  has multiplicity at least 2. Then there exist  $u_1, u_2, u_3, u_4 \in U$  such that  $a_1 + u_1 = a_2 + u_2$  and  $a_1 + u_3 = a_2 + u_4$ . Then we obtain  $u_1 + u_4 = u_2 + u_3$ , which contradicts the fact that  $U$  is a Sidon set.  $\square$

**Proposition 1.6.** *Let  $U \subseteq [n]$  be a Sidon set, and  $I \subseteq [n] - U$  such that  $U \cup I$  is also a Sidon set. Then  $I$  is an independent set in  $H^U(A)$ .*

*Proof.* Otherwise, there exist  $a_1, a_2 \in I$  and  $u_1, u_2 \in U$  such that  $a_1 + u_1 = a_2 + u_2$ , which contradicts the fact that  $U \cup I$  is a Sidon set.  $\square$

Let  $U$  be a Sidon set of size  $\frac{\sqrt{n}}{\log n}$ . The reason to choose this size is that later in the counting process, we need to make sure that the number of choices for  $U$  is at most  $2^{O(\sqrt{n})}$ . Then by Proposition 1.6, it suffices to count the number of independent sets in  $H^U([n] - U)$ .

First, let us try to apply the standard one phase graph container method. Fix an arbitrary set  $I \subseteq A$  such that  $U \cup I$  is a Sidon set. We start with  $A = [n] - U$  and  $T = \emptyset$ . In each iteration, we take a vertex  $v$  of maximum degree in  $H^U(A)$ . If  $v \in I$ , then we let  $T = T \cup \{v\}$ , and  $A = A - \{v\} - N(v)$ . Otherwise, let  $T = T$ , and  $A = A - \{v\}$ . We stop the algorithm when  $A$  is small enough, i.e.  $A = O(\sqrt{n})$ .

Note that this method works only if  $H^U(A)$  has many edges. By the supersaturation, i.e. Lemma 1.4, we would require  $|A||U| \geq 6n$ , that is,  $|A| \geq \sqrt{n} \log n$ . Therefore, the size of  $A$  cannot reach  $O(\sqrt{n})$ . Moreover, it turns out that we can prove  $|T| \leq \frac{\sqrt{n}}{\log n}$  is true only when  $|A| \geq \sqrt{n} \log^3 n$ . This is because that when  $A$  is small, the maximum degree of  $H^U(A)$  also becomes small, and then we need add more vertices in  $T$  to shrink the set  $A$ . The standard one phase graph container method fails to work.

*Proof of Theorem 1.3 [7].* First, we shall apply a multi-phase graph container method to get the certificates. Fix a Sidon set  $I$ . Let  $U$  be an arbitrary subset of  $I$  of size  $\frac{\sqrt{n}}{\log n}$ .

**Phase I.** We start with  $A = [n] - U$  and  $T = \emptyset$ . We shall apply the graph container algorithm on the graph  $H^U(A)$  in the range  $|A| \geq 6\sqrt{n} \log^3 n$ .

To be more precise, in each iteration, we take a vertex  $v \in A$  of maximum degree in  $H^U(A)$ . If  $v \in I$ , then we let  $T = T \cup \{v\}$ , and  $A = A - \{v\} - N(v)$ . Otherwise, let  $T = T$ , and  $A = A - \{v\}$ . We stop the algorithm when the size of  $A$  reaches  $6\sqrt{n} \log^3 n$ .

In the end of Phase I, let  $T_I := T$ , and  $A_I := A$ . By Lemma 1.4 and  $|A| \geq 6\sqrt{n} \log^3 n$ , in each iteration, we obtain that the average degree of  $H^U(A)$  is

$$d(H^U(A)) = \frac{2e(H^U(A))}{|A|} \geq \frac{|A||U|^2}{6n} \geq |U| \log^2 n.$$

Therefore, we have  $|T_I| \leq \frac{n}{|U| \log^2 n} = \frac{\sqrt{n}}{\log n}$ .

**Phase II.** We start with  $A = A_I$  and  $T = \emptyset$ . We shall apply the graph container algorithm on the graph  $H^U(A)$  in the range  $6\sqrt{n} \log^3 n > |A| \geq 6\sqrt{n} \log n$ .

In the end of Phase II, let  $T_{II} := T$ , and  $A_{II} := A$ . Note that in this phase, we still keep  $|A||U| \geq 6n$ , and then  $d(H^U(A)) \geq \frac{|A||U|^2}{6n} \geq |U|$ . Therefore,  $|T_{II}| \leq \frac{6\sqrt{n} \log^3 n}{|U|} = 6 \log^4 n$ .

**Phase III<sub>j</sub>**, for  $1 \leq j \leq \log \log n$ . For  $j = 1$ , we start with  $A = A_{II}$  and  $T = \emptyset$ ; for  $j \geq 2$ , we start with  $A = A_{j-1}$  and  $T = \emptyset$ . Let  $U_j$  be an arbitrary subset of  $A \cap I$  of size  $2^j \frac{\sqrt{n}}{\log n}$ . We shall apply the graph container algorithm on the graph  $H^{U_j}(A)$  in the range  $\frac{6\sqrt{n} \log n}{2^{j-1}} > |A| \geq \frac{6\sqrt{n} \log n}{2^j}$ .

In the end of Phase III<sub>j</sub>, let  $T_j := T$ , and  $A_j := A$ . Note that in this phase, we still keep  $|A||U_j| \geq 6n$ , and similarly, we have  $|T_j| \leq \frac{6\sqrt{n} \log n}{2^{j-1}} / 2^j \frac{\sqrt{n}}{\log n} \leq 3 \log^2 n$ .

From the multi-phase graph container algorithm, for each Sidon set  $I$ , we build a certificate  $\{U, U_1, \dots, U_{\log \log n}\} \cup \{T_I, T_{II}, T_1, \dots, T_{\log \log n}\}$ , which uniquely determines a set sequence  $\{A_I, A_{II}, A_1, \dots, A_{\log \log n}\}$ . Let  $C = A_{\log \log n}$ . Then the following conditions are satisfied:

- $U = \frac{\sqrt{n}}{\log n}$ ;
- $U_j = 2^j \frac{\sqrt{n}}{\log n}$  and  $U_j \subseteq A_{j-1}$ ;
- $T_I \subseteq [n]$ ,  $|T_{II}| \leq \log^4 n$ , and  $|T_j| \leq 3 \log^2 n$  for all  $j \geq 1$ ;
- $I \subseteq C \cup T_I \cup T_{II} \cup \bigcup_{j=1}^{\log \log n} T_j$ , and  $|C| = O(\sqrt{n})$ .

Note that the third condition indicates that the number of choices for  $\{T_I, T_{II}, T_1, \dots, T_{\log \log n}\}$  is  $2^{O(\sqrt{n})}$ .

**Claim 1.7.**

$$\prod_{j=1}^{\log \log n} \binom{\frac{6\sqrt{n} \log n}{2^{j-1}}}{2^j \frac{\sqrt{n}}{\log n}} = 2^{O(\sqrt{n})}.$$

*Proof.* Let  $x = \frac{6\sqrt{n}\log n}{2^{\log\log n-1}} = 12\sqrt{n}$ . Then the left side is equal to

$$\begin{aligned} \prod_{i=0}^{\log\log n-1} \binom{2^i x}{\frac{12n}{2^{2^i}}} &\leq \prod_{i=0}^{\log\log n-1} \left( \frac{e \cdot x 2^{2^i}}{12n} \right)^{\frac{12n}{2^{2^i}}} = \prod_{i=0}^{\log\log n-1} (12e 2^{2^i})^{\frac{\sqrt{n}}{2^i}} \\ &\leq \prod_{i=0}^{\log\log n-1} 2^{\lceil (\log(12e)+2i)2^{-i} \rceil \sqrt{n}} \leq 2^{\lceil \sum_{i=0}^{\infty} (\log(12e)+2i)2^{-i} \rceil \sqrt{n}} \\ &= 2^{(2\log(12e)+4)\sqrt{n}}, \end{aligned}$$

where the first inequality follows from the Stirling's formula. □

Therefore, the number of Sidon sets is at most

$$2^{|C|} 2^{O(\sqrt{n})} \binom{n}{\frac{\sqrt{n}}{\log n}} \prod_{j=1}^{\log\log n} \binom{|A_{j-1}|}{2^j \frac{\sqrt{n}}{\log n}} = 2^{O(\sqrt{n})} \prod_{j=1}^{\log\log n} \binom{\frac{6\sqrt{n}\log n}{2^{j-1}}}{2^j \frac{\sqrt{n}}{\log n}} = 2^{O(\sqrt{n})}.$$

□

## References

- [1] W. G. Brown. *On the non-existence of a type of regular graphs of girth 5*. Canadian Journal of Mathematics 19 (1967): 644–648.
- [2] P. J. Cameron and P. Erdős. *On the number of sets of integers with various properties*. Number theory (Banff, AB, 1988), 61–79. (1990)
- [3] S. Chowla. *Solution of a problem of Erdős and Turán in additive-number theory*. Proc. Nat. Acad. Sci. India. Sect. A. 14 (1944), 1–2. 1, 1.2, 7.2.
- [4] P. Erdős. *On a Problem of Sidon in Additive Number Theory and on Some Related Problems Addendum*. Journal of the London Mathematical Society 19, no. 76 Part 4 (1944): 208–208.
- [5] P. Erdős, A. Rényi, and Vera T. Sós. *On a problem of graph theory*. Publ. Math. Inst. Hungar. Acad. Sci 7 (1962): 215–235.
- [6] P. Erdős, P. Turán. *On a problem of Sidon in additive number theory, and on some related problems*. Journal of the London Mathematical Society 1, no. 4 (1941): 212–215.

- [7] Y. Kohayakawa, S. J. Lee, V. Rödl, W. Samotij. *The number of Sidon sets and the maximum size of Sidon sets contained in a sparse random set of integers*. Random Structures & Algorithms 46, no. 1 (2015): 1–25.
- [8] I. Reiman. *Über ein Problem von K. Zarankiewicz*. Acta Mathematica Hungarica 9.3–4 (1958): 269–273.
- [9] D. Saxton and A.G. Thomason, *Hypergraph Containers*. Inventiones Mathematicae, 201 (2015), 925–992.
- [10] J. Singer. *A theorem in finite projective geometry and some applications to number theory*. Transactions of the American Mathematical Society 43, no. 3 (1938): 377–385.