

A PROBLEM OF RAMANUJAN, ERDŐS AND KÁTAI ON THE ITERATED DIVISOR FUNCTION

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ABSTRACT. We determine asymptotically the maximal order of $\log d(d(n))$, where $d(n)$ is the number of positive divisors of n . This solves a problem first put forth by Ramanujan in 1915.

1 Introduction

Let $d(n)$ denote the number of positive divisors of an integer n . The extreme large values of $d(n)$ were studied by Wigert [10], (see also [4, Theorem 432]). Wigert proved that

$$m_1(x) := \max_{n \leq x} \log d(n) \sim (\log 2) \frac{\log x}{\log_2 x}.$$

Here $\log_k x$ denotes the k -th iterate of the logarithm. The lower bound comes from considering integers of the form $N_k = p_1 \cdots p_k$, where p_j denotes the j th smallest prime. Here $d(N_k) = 2^k$, while $\log N_k \sim k \log k$ by the prime number theorem. In his seminal 1915 paper on highly composite numbers [7], Ramanujan gave a more precise asymptotic for $m_1(x)$. At the very end of his paper, Ramanujan posed the problem of finding the extreme large values of $d(d(n))$. By considering integers of the form

$$(1.1) \quad 2^1 \cdot 3^2 \cdot 5^4 \cdots p_k^{p_k-1},$$

Ramanujan showed that

$$m_2(x) := \max_{n \leq x} \log d(d(n)) \geq (\sqrt{2} \log 4 + o(1)) \frac{\sqrt{\log x}}{\log_2 x}.$$

The problem of finding the order of $m_2(x)$ has been mentioned in Erdős [1], Ivić [5], and has been mentioned by Ivić in problem sessions in Ottawa [6] and Oberwolfach.

Erdős and Kátai [3] showed $m_2(x) = (\log x)^{1/2} (\log_2 x)^{O(1)}$ (see (4.1) on p. 270 of [3]). Twenty years later Erdős and Ivić [2] improved the upper bound to

$$m_2(x) \ll \left(\frac{\log x \log_2 x}{\log_3 x} \right)^{1/2}.$$

Smati [8, 9] gave a further improvement

$$m_2(x) \ll \sqrt{\log x},$$

the best estimate known to date. Constructions similar to Ramanujan's seem rather natural, and one might expect that $m_2(x) \ll \frac{\sqrt{\log x}}{\log_2 x}$. This is indeed the case, as we now show. More precisely, we prove an asymptotic formula for $m_2(x)$ with an error term.

Theorem 1. *We have*

$$m_2(x) = \frac{\sqrt{\log x}}{\log_2 x} \left(c + O \left(\frac{\log_3 x}{\log_2 x} \right) \right),$$

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where

$$c = \left(8 \sum_{j=1}^{\infty} \log^2(1 + 1/j) \right)^{1/2} = 2.7959802335 \dots$$

In particular, Theorem 1 implies that

$$\limsup_{n \rightarrow \infty} \frac{\log d(d(n)) \log_2 n}{\sqrt{\log n}} = c.$$

Ramanujan's examples (1.1) are seen to be suboptimal with respect to the constant c , since $\sqrt{2} \log 4 = 1.9605 \dots$

There is a closely related problem, to estimate the extreme values of $\omega(d(n))$, where $\omega(n)$ is the number of distinct prime factors of n . In fact, both Erdős and Ivić [2] and Smati [9] obtained upper bounds for $d(d(n))$ by first bounding $\omega(d(n))$ and then using the elementary inequality $\log d(m) \ll (\log_2 m) \omega(m)$ (see, e.g., Lemme 3.3 of [8] or Lemma 3.2 below). For this problem, Ramanujan's examples (1.1) are essentially optimal, providing the true order and constant in the asymptotic for $w(x) = \max_{n \leq x} \omega(d(n))$.

Theorem 2. *We have*

$$w(x) = \frac{\sqrt{\log x}}{\log_2 x} \left(\sqrt{8} + O\left(\frac{\log_3 x}{\log_2 x}\right) \right),$$

Previously, Erdős and Ivić [2] had shown

$$w(x) \ll \left(\frac{\log x \log_3 x}{\log_2 x} \right)^{1/2},$$

and later Smati [8] found the true order $w(x) \ll \frac{\sqrt{\log x}}{\log_2 x}$.

2 The lower bound in Theorem 1

Notation and basic prime number estimates. Throughout, we make use of the asymptotic

$$(2.1) \quad p_j = j(\log j + \log_2 j + O(1)),$$

which is a simple consequence of the prime number theorem with error term $\pi(x) = \frac{x}{\log x} + O\left(\frac{x}{\log^2 x}\right)$. Here $\pi(x)$ is the number of primes which are $\leq x$. We also denote by $\Omega(n)$ the number of prime power divisors of n .

Proof of the lower bound in Theorem 1. Let x be large and define $\varepsilon = 10 \frac{\log_3 x}{\log_2 x}$. Let

$$(2.2) \quad t = \left\lfloor \left(\frac{8 \log 2}{c} - \varepsilon \right) \frac{\sqrt{\log x}}{\log_2 x} \right\rfloor, \quad a_i = \left\lfloor \frac{1}{2^{i/t} - 1} \right\rfloor \quad (1 \leq i \leq t),$$

and let

$$n = (p_1 \cdots p_{a_1})^{p_1 - 1} (p_{a_1 + 1} \cdots p_{a_1 + a_2})^{p_2 - 1} \cdots (p_{a_1 + \cdots + a_{t-1} + 1} \cdots p_{a_1 + \cdots + a_t})^{p_t - 1}.$$

The Taylor expansion of $\exp\left(\frac{\log 2}{t}\right)$ shows that $a_1 = \lfloor (2^{1/t} - 1)^{-1} \rfloor = t / \log 2 + O(1)$. By (2.2), for every positive integer j , there are $y_j := \lfloor \frac{\log(1+1/j)t}{\log 2} \rfloor$ indices i with $a_i \geq j$. Also, $a_1 + \cdots + a_t \ll t \log t$. Using (2.1), we have $\log p_{a_1 + \cdots + a_i} \leq \log t + 2 \log_2 t + O(1)$, hence

$$\log n \leq \sum_{i=1}^t a_i (p_i - 1) \log p_{a_1 + \cdots + a_i} \leq (\log^2 t + 3(\log_2 t) \log t + O(\log t)) \sum_{i=1}^t i a_i.$$

From $y_j = O(t/j)$ and the definition of c we obtain

$$(2.3) \quad \begin{aligned} \sum_{i=1}^t i a_i &= \sum_{j \leq a_1} \frac{y_j(y_j + 1)}{2} = \frac{1}{2} \sum_{j=1}^{\infty} \left(\frac{\log(1 + 1/j)}{\log 2} \right)^2 t^2 + O(t \log t) \\ &= \frac{c^2}{16(\log 2)^2} t^2 + O(t \log t). \end{aligned}$$

From the definition of t , $\log t = \frac{1}{2} \log_2 x - \log_3 x + O(1)$ and $\log_2 t = \log_3 x + O(1)$. Thus,

$$\log n \leq \left(1 + \frac{2 \log_3 x + O(1)}{\log_2 x} \right) \left(1 - \frac{c\varepsilon}{8 \log 2} \right)^2 \log x.$$

Hence, if x is large enough, then $n \leq x$. From the definition of n above, we have $d(n) = p_1^{\alpha_1} \cdots p_t^{\alpha_t}$. Therefore,

$$(2.4) \quad \begin{aligned} \log m_2(x) &\geq \log d(d(n)) = \sum_{i=1}^t \log(a_i + 1) = \sum_{j \geq 1} (y_j - y_{j+1}) \log(j + 1) = \sum_{j \geq 1} y_j \log(1 + 1/j) \\ &= \sum_{j \leq a_1} \left(\frac{\log^2(1 + 1/j)}{\log 2} t + O(1/j) \right) \\ &= \frac{c^2}{8 \log 2} t + O(\log t) \\ &= \frac{\sqrt{\log x}}{\log_2 x} \left(c + O\left(\frac{\log_3 x}{\log_2 x} \right) \right). \end{aligned}$$

□

3 Proof of the upper bound in Theorem 1

Lemma 3.1. *Let $m_N = \min\{m : d(m) = N\}$ and write $m_N = p_1^{\alpha_1} \cdots p_r^{\alpha_r}$. We have*

- (i) $\alpha_1 \geq \cdots \geq \alpha_r$,
- (ii) $N' | N$ implies $m_{N'} \leq m_N$,
- (iii) for each integer $k \geq 1$, if $p_j > p_{r+1}^{1/2^k}$, then $\Omega(\alpha_j + 1) \leq k$.

Remark 1. Using (2.1) and taking $k = 1$, we see from (iii) that if r is large, then $\alpha_j + 1$ is prime for $\sqrt{r} < j \leq r$. Also, by (iii), $\Omega(\alpha_j + 1) \ll \log_2 r$ for all j .

Proof. (i) This is trivial and was observed by Ramanujan [7, (32)].

(ii) If $N' | N$, we can find $\alpha'_j \leq \alpha_j$ for each j such that $N' = (\alpha'_1 + 1) \cdots (\alpha'_r + 1)$, and clearly $m_{N'} \leq p_1^{\alpha'_1} \cdots p_r^{\alpha'_r} \leq m_N$.

(iii) If $p_j > p_{r+1}^{1/2^k}$ and $\Omega(\alpha_j + 1) > k$, then there are integers a, b with $\alpha_j + 1 = ab$, $a \geq 2$ and $b \geq 2^k$. Letting

$$m^* = p_j^{b-1} p_{r+1}^{a-1} \prod_{i \neq j} p_i^{\alpha_i},$$

we see that $d(m^*) = d(m_N) = N$, but

$$\frac{m^*}{m_N} = p_j^{b-1-\alpha_j} p_{r+1}^{a-1} = (p_j^{-b} p_{r+1})^{a-1} < 1,$$

a contradiction. □

Lemma 3.2. For every $\varepsilon > 0$, and for $\omega(n) = s \geq 2$ we have

$$d(n) \ll_{\varepsilon} \left(\frac{(2 + \varepsilon) \log n}{s \log s} \right)^s.$$

Proof. Write the prime factorization of n as $n = q_1^{a_1} \cdots q_s^{a_s}$, where $q_1 < \cdots < q_s$. Using the arithmetic mean - geometric mean inequality and that $q_i \geq p_i$, we have

$$d(n) \leq \prod_{i=1}^s (2a_i) \leq 2^s \prod_{i=1}^s (a_i \log q_i) \prod_{i=1}^s (\log p_i)^{-1} \leq \left(\frac{2 \log n}{s} \right)^s \frac{(\log s)^{\pi(s)-s}}{\log 2},$$

the last inequality coming from excluding factors corresponding to $3 \leq p_i < s$. Finally, the prime number theorem implies $(\log s)^{\pi(s)} \leq (\log s)^{O(s/\log s)} \ll_{\varepsilon} (1 + \varepsilon/2)^s$. \square

Remark. Lemma 3.2 is fairly sharp. For example, from the inequality $s = \omega(n) \leq (1 + o(1)) \frac{\log n}{\log 2}$, and the observation that $m_1(x)$ is monotonically increasing, we immediately obtain Wigert's upper bound for $\log d(n)$.

The following is the key lemma, which explains the constant c .

Lemma 3.3. Let a_1, \dots, a_t be positive integers.

(a) we have

$$\sum_{i=1}^t \log(a_i + 1) \leq \frac{c}{2} \left(\sum_{i=1}^t ia_i \right)^{1/2}.$$

Moreover, the constant $c/2$ is best possible.

(b) If $a_i \geq A$ for all i , where A is a positive integer, then

$$\sum_{i=1}^t \log(a_i + 1) \leq \left(\frac{1 + \log^2(A + 1)}{A} \sum_{i=1}^t ia_i \right)^{1/2}.$$

Proof. (a) Without loss of generality, suppose $a_1 \geq \cdots \geq a_t$. Let $y_j = \#\{i : a_i \geq j\}$. Then

$$\sum_{i=1}^t ia_i = \sum_{j \geq 1} \frac{y_j(y_j + 1)}{2} \geq \frac{1}{2} \sum_{j \geq 1} y_j^2.$$

By partial summation and the Cauchy-Schwarz inequality,

$$\begin{aligned} \sum_{i=1}^t \log(a_i + 1) &= \sum_{j \geq 1} (y_j - y_{j+1}) \log(j + 1) = \sum_{j \geq 1} y_j \log(1 + 1/j) \\ (3.1) \qquad \qquad \qquad &\leq \left(\sum_{j \geq 1} y_j^2 \right)^{1/2} \left(\frac{c^2}{8} \right)^{1/2}. \end{aligned}$$

Moreover, the inequality in (3.1) is an equality if and only if for some real Y , $y_j = Y \log(1 + 1/j)$ for every j . As the y_j are integers, this cannot happen. However, we can come very close to equality in (3.1) by taking t large and choosing the a_i by (2.2), so that $y_j = \lfloor \frac{\log(1+1/j)}{\log 2} t \rfloor$. By (2.3) and (2.4), we have in this case

$$\sum_{i=1}^t \log(a_i + 1) = \frac{c^2}{8 \log 2} t + O(\log t), \quad \sum_{i=1}^t ia_i = \frac{c^2}{16(\log 2)^2} t^2 + O(t \log t),$$

whence

$$\sum_{i=1}^t \log(a_i + 1) = \frac{c}{2} \left(1 + O\left(\frac{\log t}{t}\right) \right) \left(\sum_{i=1}^t ia_i \right)^{1/2}.$$

(b) Observe that $y_1 = y_2 = \cdots = y_A$. Arguing similarly to (3.1), we obtain

$$\begin{aligned} \sum_{i=1}^t \log(a_i + 1) &= \frac{\log(A+1)}{A} (y_1 + \cdots + y_A) + \sum_{j>A} y_j \log(1 + 1/j) \\ &\leq \left(\sum_{j \geq 1} y_j^2 \right)^{1/2} \left(A \left(\frac{\log(A+1)}{A} \right)^2 + \sum_{j>A} \log^2(1 + 1/j) \right)^{1/2}. \end{aligned}$$

Observing that $\log(1 + 1/j) < 1/j$ and $\sum_{j>A} 1/j^2 < 1/A$, we obtain (b). \square

The next lemma is trivial.

Lemma 3.4. *For any positive integer m , $m \geq \sum_{p|m} p$.*

Proof of Theorem 1, upper bound. Let n be large, let $N = d(n)$ and factor $N = N'N''$, where

$$N' = u_1^{b_1} \cdots u_w^{b_w}, \quad N'' = q_1^{a_1} \cdots q_s^{a_s},$$

where $u_1 < \cdots < u_w$, $q_1 < \cdots < q_s$ are primes, $b_i > (\log_2 n)^6$ for every i and $a_i \leq (\log_2 n)^6$ for every i .

Write $m_{N'} = p_1^{\beta_1} \cdots p_h^{\beta_h}$. By Lemma 3.1 (ii), $m_{N'} \leq m_N \leq n$, so that $h \ll \log n$. By Lemma 3.1 (iii), $\Omega(\beta_i + 1) \ll \log_2 h \ll \log_3 n$ for every i . Since $d(m_{N'}) = (\beta_1 + 1) \cdots (\beta_h + 1) = N'$, for each $j \leq h$ there are $\gg \frac{b_j}{\log_3 n}$ values of i for which $u_j | (\beta_i + 1)$. Thus, by Lemma 3.4,

$$\begin{aligned} \log n \geq \log m_{N'} &\geq (\log 2) \sum_{i=1}^h \beta_i \geq \frac{\log 2}{2} \sum_{i=1}^h (\beta_i + 1) \\ &\geq \frac{\log 2}{2} \sum_{i=1}^h \sum_{p | (\beta_i + 1)} p \gg \sum_{j=1}^w u_j \frac{b_j}{\log_3 n} \geq \frac{1}{\log_3 n} \sum_{j=1}^w j b_j. \end{aligned}$$

Combining this estimate with Lemma 3.3 (b) with $A = (\log_2 n)^6$ gives

$$(3.2) \quad \log d(N') = \sum_{j=1}^w \log(b_j + 1) \ll \frac{\log_3 n}{(\log_2 n)^3} \left(\sum_{j=1}^w j b_j \right)^{1/2} \ll \frac{(\log n)^{1/2} (\log_3 n)^{3/2}}{(\log_2 n)^3}.$$

Next, we bound $d(N'')$.

Case 1) If $s \leq \frac{(\log n)^{1/2}}{(\log_2 n)^3}$, Lemma 3.2 implies that $\log d(N'') \ll \frac{(\log n)^{1/2}}{(\log_2 n)^2}$.

Case 2) Now suppose that $s > \frac{(\log n)^{1/2}}{(\log_2 n)^3}$. Write $m_{N''} = p_1^{\alpha_1} \cdots p_r^{\alpha_r}$. By Lemma 3.1 (iii),

$$(3.3) \quad r \leq \Omega(N'') = \sum_{j=1}^s a_j = \sum_{i=1}^r \Omega(\alpha_i + 1) \leq r + \sum_{k \geq 2} \pi(p_{r+1}^{1/2^k}) = r + O((r/\log r)^{1/2}).$$

In particular, $r + O((r/\log r)^{1/2}) \geq a_1 + \cdots + a_s \geq s$, so $r \gg s > \frac{(\log n)^{1/2}}{(\log_2 n)^3}$. Thus, for large enough n , $a_1 + \cdots + a_s \leq r + \sqrt{r}$. Also by Lemma 3.1 (iii), $\alpha_j + 1$ is prime for $j > \sqrt{r}$. Let $\varepsilon = 20 \frac{\log_3 n}{\log_2 n}$. By the lower bound on s , and using $a_i \leq (\log_2 n)^6$,

$$(3.4) \quad \sum_{j > s - s^{1-\varepsilon}} a_j \geq s^{1-\varepsilon} \geq 2 (s (\log_2 n)^6)^{1/2} \geq 2 (\Omega(N''))^{1/2} \geq 2\sqrt{r},$$

hence, using (3.3),

$$\sum_{j \leq s - s^{1-\varepsilon}} a_j \leq \Omega(N'') - 2\sqrt{r} \leq r - \sqrt{r}.$$

Using Lemma 3.1 (i), $\alpha_i + 1 = q_1$ for $r - a_1 < i \leq r$, and similarly for each $j \leq s - s^{1-\varepsilon}$, $\alpha_i + 1 = q_j$ for $r - (a_1 + \cdots + a_j) < i \leq r - (a_1 + \cdots + a_{j-1})$. We obtain

$$\begin{aligned} \log m_{N''} &\geq \sum_{\sqrt{r} < i \leq r} \alpha_i \log p_i \geq \sum_{j \leq s - s^{1-\varepsilon}} (q_j - 1) \sum_{m=r-(a_1+\cdots+a_j)+1}^{r-(a_1+\cdots+a_{j-1})} \log p_m \\ &\geq \sum_{j \leq s - s^{1-\varepsilon}} (p_j - 1) a_j \log(r - (a_1 + \cdots + a_j)). \end{aligned}$$

By (3.4), uniformly for $j \leq s - s^{1-\varepsilon}$ we have

$$r - (a_1 + \cdots + a_j) = r - \Omega(N'') + a_{j+1} + \cdots + a_s \geq s - j - \sqrt{r} \geq \frac{1}{2} s^{1-\varepsilon}.$$

Using (2.1), $p_j \geq j \log j + 1$ for large j . Hence, by Lemma 3.1 (ii),

$$\begin{aligned} \log n &\geq \log m_{N''} \geq \sum_{s^{1-\varepsilon} \leq j \leq s - s^{1-\varepsilon}} (j \log j) a_j (\log s + O(\log_3 n)) \\ &\geq (1 + O(\varepsilon)) \frac{(\log_2 n)^2}{4} \sum_{s^{1-\varepsilon} \leq j \leq s - s^{1-\varepsilon}} j a_j. \end{aligned}$$

By the definition of ε , $s^\varepsilon \gg (\log_2 n)^9$. Also, trivially $\sum_{j=1}^s j a_j \geq 1 + 2 + \cdots + s \geq \frac{1}{2} s^2$. Recalling that $a_j \leq (\log_2 n)^6$ for every j , we have

$$\begin{aligned} \sum_{s^{1-\varepsilon} \leq j \leq s - s^{1-\varepsilon}} j a_j &= \sum_{j=1}^s j a_j + O(s^{2-\varepsilon} (\log_2 n)^6) = \sum_{j=1}^s j a_j + O(s^2 (\log_2 n)^{-3}) \\ &= (1 + O(1/\log_2 n)) \sum_{j=1}^s j a_j. \end{aligned}$$

Combining the last two inequalities gives

$$\log n \geq \left(1 + O\left(\frac{\log_3 n}{\log_2 n}\right)\right) \frac{(\log_2 n)^2}{4} \sum_{j=1}^s j a_j.$$

Applying Lemma 3.3 (a), we conclude that

$$(3.5) \quad \log d(N'') = \sum_{j=1}^s \log(a_j + 1) \leq \frac{c}{2} \left(\sum_{j=1}^s j a_j\right)^{1/2} \leq c \frac{\sqrt{\log n}}{\log_2 n} \left(1 + O\left(\frac{\log_3 n}{\log_2 n}\right)\right).$$

Recall that we have a smaller upper bound for $\log d(N'')$ in case 1). Finally, using $d(d(n)) = d(N')d(N'')$ and combining (3.2) and (3.5), we obtain the desired upper bound for $d(d(n))$. \square

4 Proof of Theorem 2

Proof of Theorem 2. For the lower bound, let x be large and put $n = \prod_{i=1}^s p_i^{p_i-1}$, where s is the largest integer such that $n \leq x$. Recall that p_j is the j -th smallest prime. Then $d(n) = \prod_{i=1}^s p_i$, thus $\omega(d(n)) = s$. By (2.1),

$$\log n = \sum_{i=1}^s (p_i - 1) \log p_i = \sum_{i=1}^s i \log^2 i + O(i \log i \log_2 i) = \frac{1}{2} s^2 \log^2 s + O(s^2 \log s \log_2 s).$$

Solving for s gives $s = \frac{\sqrt{8 \log n}}{\log_2 n} + O\left(\frac{\sqrt{\log n \log_3 n}}{\log_2^2 n}\right)$. We now prove a lower bound on n . Since $p_{s+1} \sim s \log s \sim \sqrt{2 \log n} \ll \sqrt{\log x}$ by (2.1), we have

$$x \geq n \geq x p_{s+1}^{-p_{s+1}} = x \exp\left(-O\left(\sqrt{\log x \log_2 x}\right)\right).$$

That is, $\log n = \log x + O(\sqrt{\log x \log_2 x})$. Therefore, $s = \frac{\sqrt{8 \log x}}{\log_2 x} + O\left(\frac{\sqrt{\log x \log_3 x}}{\log_2^2 x}\right)$.

Now let n be a large, positive integer factored as $n = n_1 n_2$, $n_1 = \prod_{i=1}^r q_i^{a_i}$, $n_2 = \prod_{i=1}^{r'} (q'_i)^{a'_i}$, where q_i, q'_i are primes, $q_i > P$ and $q'_i \leq P$ for each i , where $P = \frac{\sqrt{\log n}}{\log_2 n}$. We have

$$(4.1) \quad \omega(d(n)) \leq \omega(d(n_1)) + \omega(d(n_2)).$$

Since $\omega(n_2) \leq \pi(P) \ll \frac{\sqrt{\log n}}{(\log_2 n)^2}$, Lemma 3.2 implies $\log d(n_2) \ll \sqrt{\log n} / \log_2 n$. Applying the elementary inequality $\omega(u) \ll \frac{\log u}{\log_2 u}$ gives

$$(4.2) \quad \omega(d(n_2)) \ll \frac{\sqrt{\log n}}{(\log_2 n)^2}.$$

Next,

$$\log n_1 \geq (\log P) \sum_{i=1}^r a_i = \left(\frac{\log_2 n}{2} - \log_3 n\right) \sum_{i=1}^r a_i.$$

Letting $s = \omega(d(n_1)) = \omega(\prod (a_i + 1))$, Lemma 3.4 implies that

$$\sum_{i=1}^r a_i \geq \sum_{i=1}^r \sum_{p|(a_i+1)} (p-1) \geq \sum_{i=1}^s (p_i - 1) \geq \sum_{i=1}^s (i \log i + O(1)) = \frac{1}{2} s^2 \log s + O(s^2).$$

Here we used the one-sided inequality $p_i \geq i \log i + O(1)$ deduced from (2.1). Thus,

$$\log n \geq \log n_1 \geq \left(\frac{1}{4} + O\left(\frac{\log_3 n}{\log_2 n}\right)\right) (\log_2 n) s^2 \log s + O(s^2 \log_2 n).$$

Consider two cases: (i) $s \leq \frac{\sqrt{\log n}}{\log_2 n}$, (ii) $s > \frac{\sqrt{\log n}}{\log_2 n}$. In case (ii), we have $\frac{\log n}{\log_2^2 n} \geq \left(\frac{1}{8} + O\left(\frac{\log_3 n}{\log_2 n}\right)\right) s^2$, and we obtain in both cases

$$\omega(d(n_1)) = s \leq \frac{\sqrt{8 \log n}}{\log_2 n} + O\left(\frac{\sqrt{\log n \log_3 n}}{\log_2^2 n}\right),$$

Combining this inequality with (4.1) and (4.2), we obtain the desired upper bound for $\omega(d(n))$. \square

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