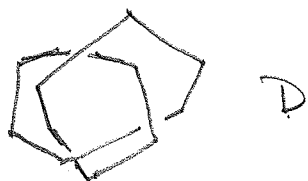


Last time we defined a knot in \mathbb{R}^3 to be a polygonal embedding $K: S^1 \rightarrow \mathbb{R}^3$ (or its image $K = K(S^1)$).

A nice projection exists and gives us a diagram of the knot



Can view \mathcal{D} as a graph with vertices of valence 2 (images of vertices of K) and vertices of valence 4 (the crossings), the latter marked w/ over/under.



Also defined equivalence relation \sim on knots (w/ 3 different descriptions).

Exercise I.1: If K, K' have the same diagram, then $K \sim K'$.

What if the diagrams are different?

Lets look at ways the diagrams could be different:

Planar isotopy:



suppose K, K' have diagrams

$\mathcal{D}, \mathcal{D}'$ in a plane \mathbb{R}^2 , say $\mathcal{D}, \mathcal{D}'$ are related by planar isotopy

if $\exists H: \mathbb{R}^2 \times [0,1] \rightarrow \mathbb{R}^2$ w/ $H_t(x) = H(x,t)$, $H_t: \mathbb{R}^2 \rightarrow \mathbb{R}^2$ homeomorphism,

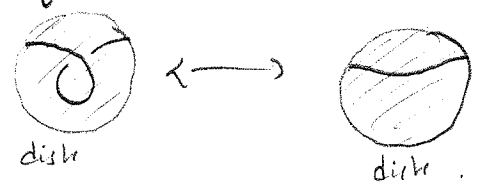
$H_0 = \text{identity on } \mathbb{R}^2$, $H_1(\mathcal{D}) = \mathcal{D}'$ respecting over/under at crossings.

then $K \sim K'$: extend H (by identifying an 3rd coordinate in \mathbb{R}^3)
 to a map $\tilde{H}: \mathbb{R}^3 \times [0,1] \rightarrow \mathbb{R}^3$, $\tilde{H}_1(K)$ is a knot with the same
 projection as K' , so $K \sim \tilde{H}_1(K) \sim K'$.

Reidemeister moves : There are 3 ways to modify the diagram
 of a knot that clearly does not change the knot type, those
 are called the Reidemeister moves ;

$D, D' \subset \mathbb{R}^2$ two diagrams, the same outside a disk ; differing
 in the disk by :

type I :



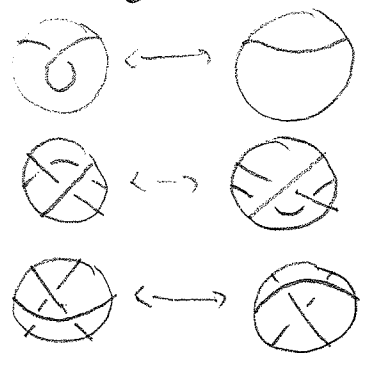
type II :



type III :



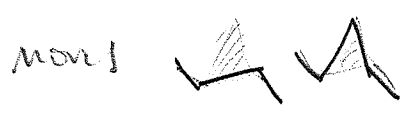
Exercise I.2 : show that
 the "other Reidemeister moves"
 below are a consequence of
 these 3 (and planar isotopy)



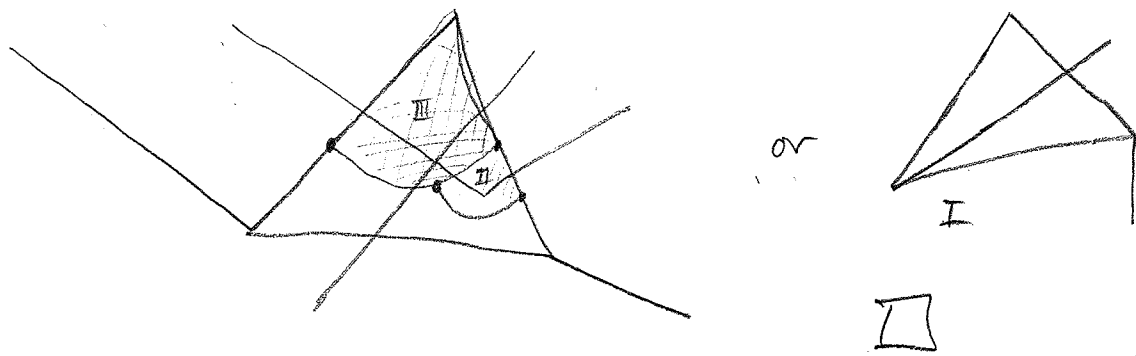
Theorem I.3, K, K' two knots are equivalent iff the diagrams
 D, D' for K, K' differ by a finite sequence of Reidemeister moves I, II, III
 and planar isotopy.

Idea of proof : Need to see if $K \sim K'$, then D, D' differ as in theorem.

Know $\exists K = K_1, K_2, \dots, K_n = K'$ w/ K_i, K_{i+1} related by triangle



By small change in projection $\pi: \mathbb{R}^3 \rightarrow \mathbb{R}^2$ we can assume π is a nice projection for every K_i . [Also note that small change of projection changes $D \& D'$ to planar isotopic diagrams]. Further, we can assume π is injective on all triangles defining moves. Now replace each triangle move with finitely many small moves of type I, II, and III. eg.



Remark: We can essentially take this theorem as defn of equivalence — i.e., $K \sim K'$ iff diagrams are related as in thm. This allows us to view knots and question of equiv. of knots as a combinatorial problem.

How does this help? [replaced existence/non existence of one type of sequence with another].

We can use this to define knot invariants:

Let $\text{Knot}(\mathbb{R}^3)$ be the set of all knots in \mathbb{R}^3 .

Defn A knot invariant is a function

$$F: \text{Knot}(\mathbb{R}^3) \rightarrow \Sigma$$

where Σ is a set with an equivalence relation \approx on it such that if $K \sim K'$, then $F(K) \approx F(K')$.

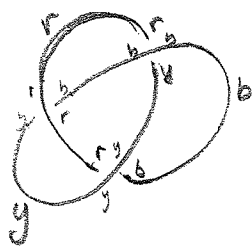
Σ could be a set of groups or abelian groups (w/ equiv. rel. \cong), or polynomials, $\mathbb{R}, \mathbb{Z}, \mathbb{C}$, w/ trivial rel. ... Anything!

How do we use Reid. moves to do this?

EX $\text{Tri}: \text{Knot}(\mathbb{R}^3) \rightarrow \mathbb{Z}$

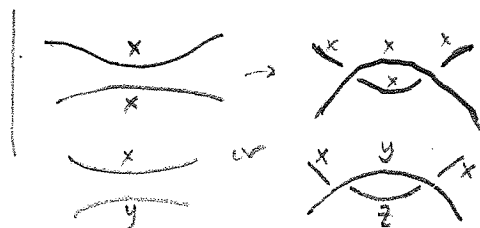
$\text{Tri}(K)$ is defined as follows: Let D be a diagram of K

$\text{Tri}(K) = \#$ of ways to color overpassing arcs of D with 3 colors (using at least 2) so that at any crossing either all 3 colors appear, or exactly one does:



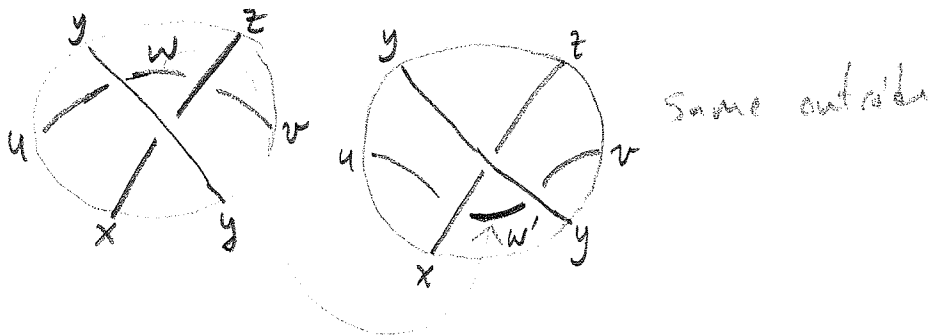
$K = \text{trefoil}$, $\text{Tri}(K) = 6$ (if two arcs are same color, all are)

If $K \sim K'$, why is $\text{Tri}(K) = \text{Tri}(K')$? check invariant under RI, II, III : that is, if $D \sim D'$ differ



by RI, II, III , then a coloring of one uniquely determines a coloring of other

III: dual-1 correspondence between arcs of D & D'



check:

- ① w, w' are determined by u, v, x, y, z — easy ($u, y \Rightarrow w$)
- ② if u, v, x, y, z, w is a coloring, then $\exists w'$ st. u, v, x, y, z, w' is a coloring (its unique by ①) — cases: $x \neq y$, exactly 3 possibilities
- ③ reverse roles of w, w'
 - $x = y$, exactly 3 possibilities.

This proves

Theorem I.4 Tri is a knot invariant

Corollary I.5 $\mathcal{D} \neq \emptyset$

Links: {A link is a collection of pairwise disjoint knots}

Defn A link is a (polygonal) embedding $L: \coprod_{i=1}^k S^1 \rightarrow \mathbb{R}^3$. We also

write $L = L(\coprod_{i=1}^k S^1)$. The restriction to one of the circles is called a component $L_1, \dots, L_k: S^1 \rightarrow \mathbb{R}^3$ (or $L_i(S^1) = L_i$) and each L_i is a knot.

Say L has k -components in this case. [a one component link is a knot?]

Same defn of equivalence, and I.1 — I.4 also hold for links.

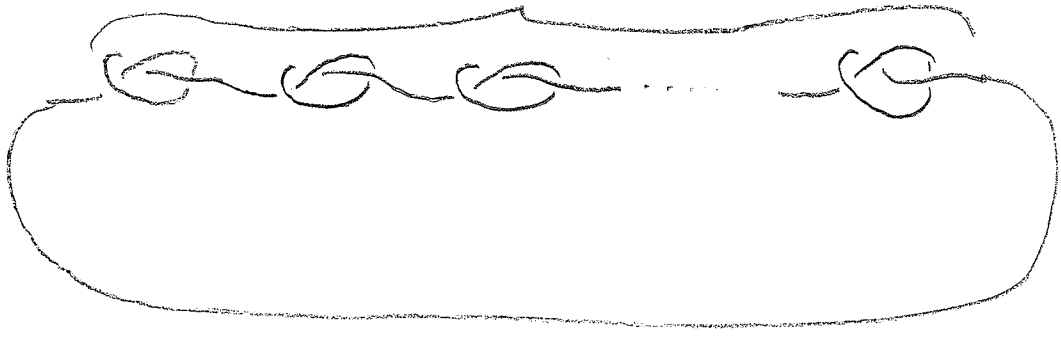
k -component unlink:



move up after Corollary I.5.

Corollary I.6: \exists infinitely many, pairwise inequivalent knots,

proof: Consider the knot K_k :
k trefoils



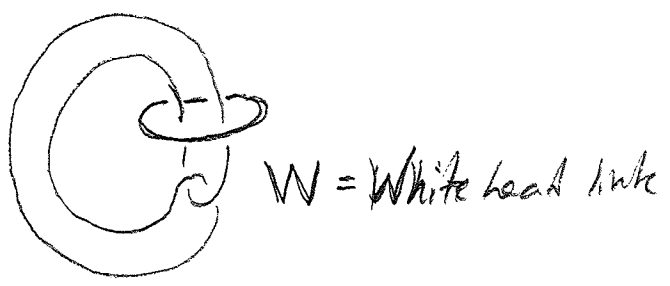
check that there are $3(3^k - 1)$ tricolorings, so

$$Tri(K_k) = 3(3^k - 1).$$

Therefore, $Tri(Knot(\mathbb{R}^3))$ is infinite, hence so is the # of equivalence classes in $Knot(\mathbb{R}^3)$.

Exercise I.3 There are only a countably infinite number of equivalence classes of knots. (of course, $Knot(\mathbb{R}^3)$ is uncountable)

Link example:



Note: $W \neq \emptyset$ since $Tri(W) = 0$, but $Tri(\emptyset) = 6$.